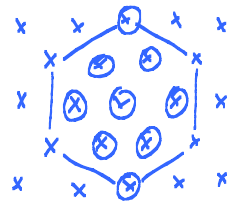


# Tutorial - Part A Outline

1. Definitions: Partition, Construction  $\text{Vol}(\Lambda)$   
Modulo  $\Lambda$
2. Figures of merit  $G(\Lambda)$
3. Dither & estimation  $\text{noise}(\Lambda)$
4. Entropy coding  $H(\Lambda)$
5. Infinite constellation  $P_e(\Lambda + \text{noise})$
6. **Asymptotic goodness**  $(n \rightarrow \infty)$
7. Error exponents
8. Nested lattices  $\Lambda_2 \subset \Lambda_1$
9. Lattice (Voronoi) shaping
10. Side-information problems  $\text{Modulo}^2(\Lambda)$
11. Gaussian networks  $\text{Modulo}^n(\Lambda)$



# Tutorial-Part A Outline

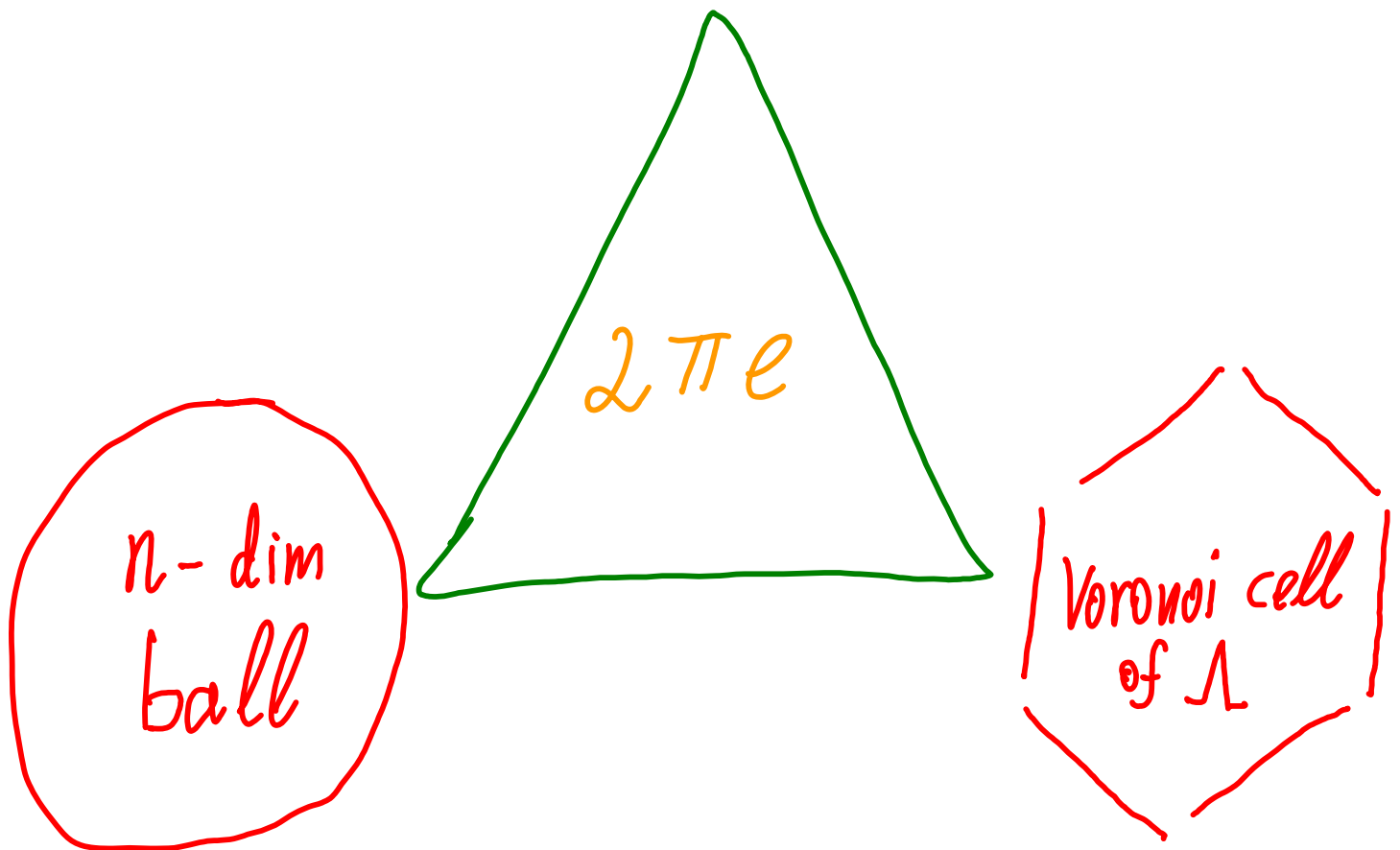
## 6. Asymptotic goodness

$$G(\Lambda_n) \xrightarrow{?} \frac{1}{2\pi e}, \text{ as } n \rightarrow \infty$$

$$\mu(\Lambda_n, p_0) \xrightarrow{?} 2\pi e, \text{ as } n \rightarrow \infty \quad \forall p_0 > 0$$

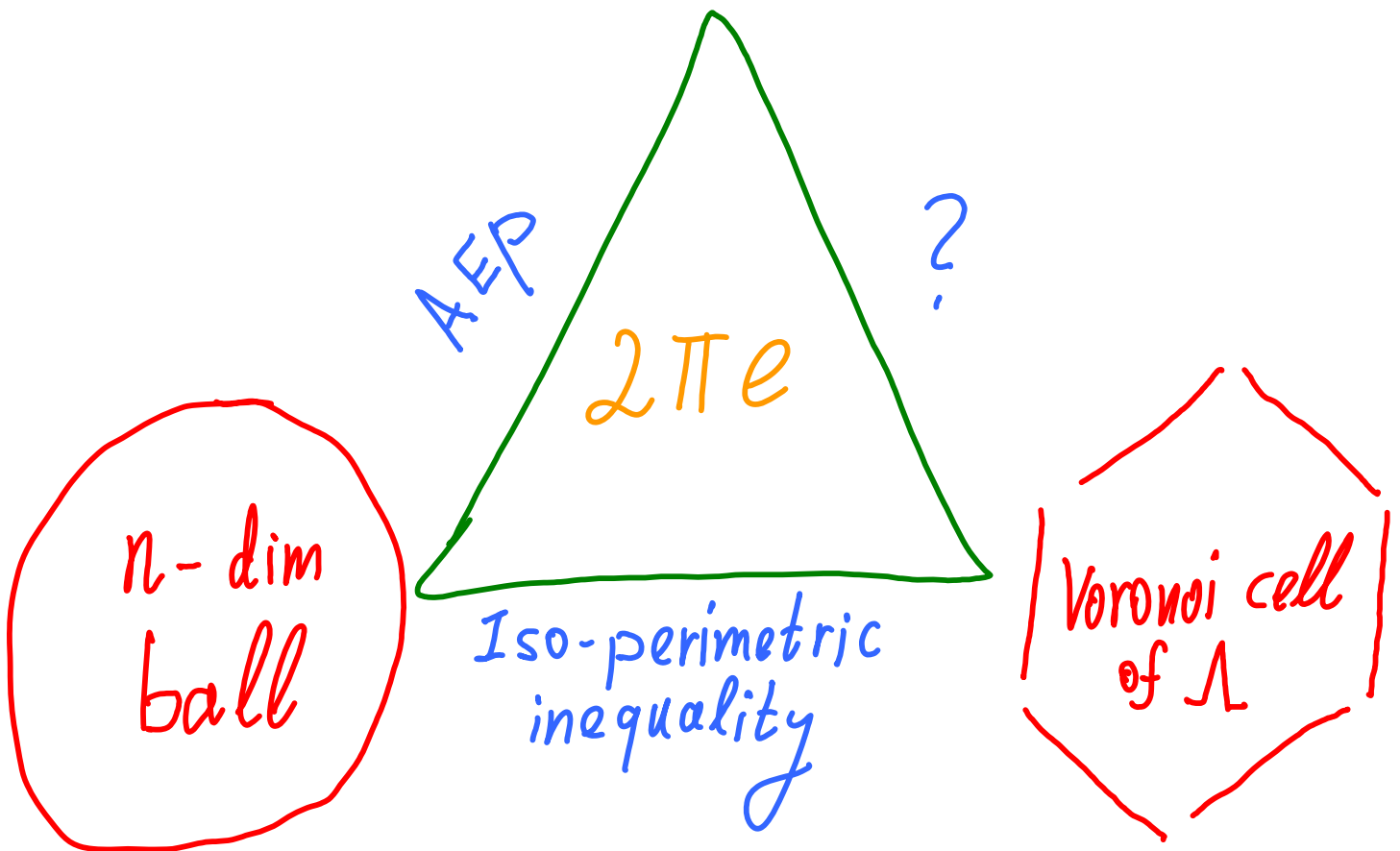
W.G.N.  $\leftrightarrow$  Ball  $\leftrightarrow$   $\Lambda$

  
white Gaussian noise



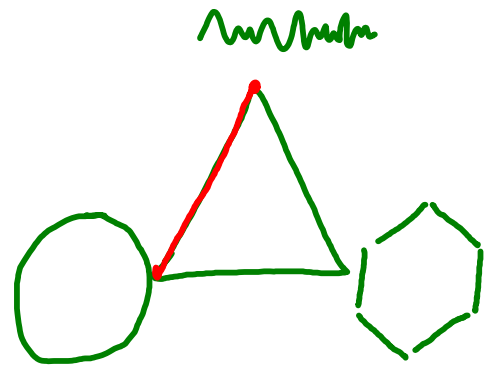
W.G.N.  $\leftrightarrow$  Ball  $\leftrightarrow$   $\Omega$

  
white Gaussian noise



# Shannon's AEP:

W.G.N.  $\rightarrow$  ball

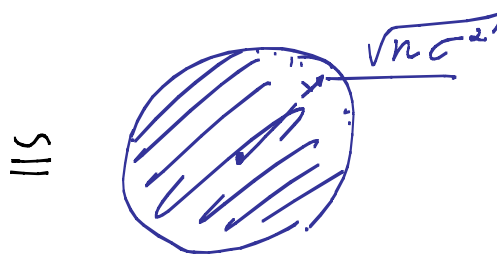
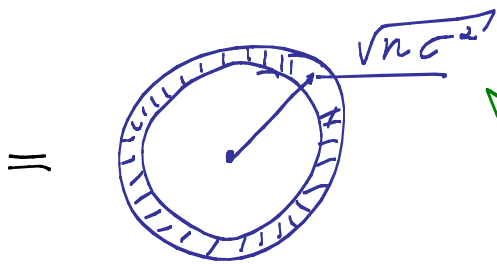


$$Z_1 \dots Z_n \sim \text{AWGN } N(0, \sigma^2)$$

$$A_\epsilon = \left\{ \mathbf{z} : \frac{1}{n} \log f_{\mathbf{z}}(\mathbf{z}) = h \pm \epsilon \right\}$$

$$= \left\{ \mathbf{z} : \|\mathbf{z}\| = \sqrt{n(\sigma^2 \pm \epsilon')} \right\}$$

AWGN  
 $f_{\mathbf{z}} \sim e^{-\frac{\|\mathbf{z}\|^2}{2\sigma^2}}$   
 $h = \frac{1}{2} \log 2\pi\sigma^2$

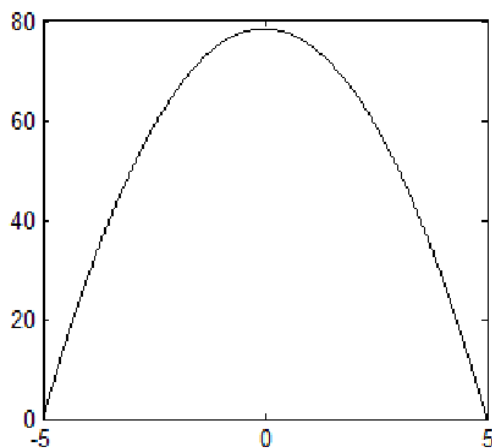
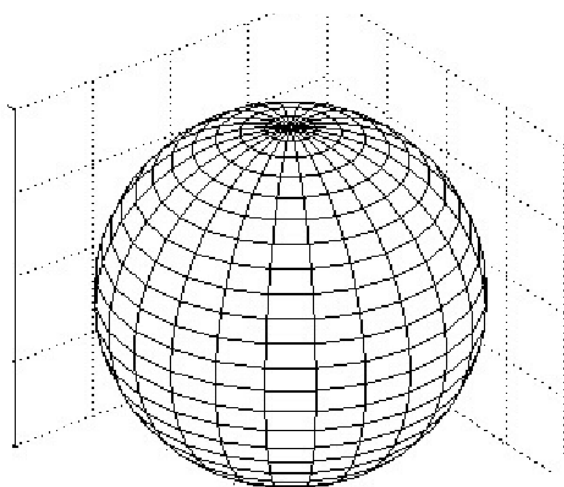
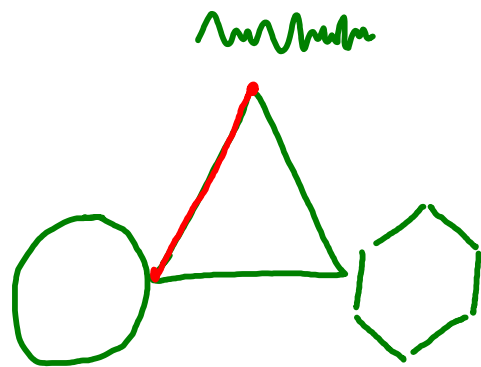


$\triangleq r_{\text{noise}}$

Thm. [AEP]: AWGN  $\approx$  Unif( $B(\mathbf{0}, \sqrt{n\sigma^2})$ )

"Reverse" AEP:

W.G.N. ← ball

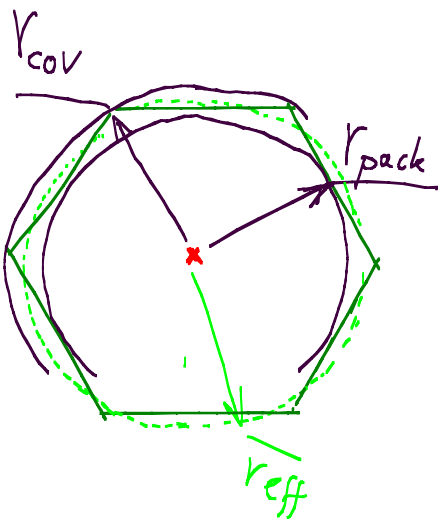
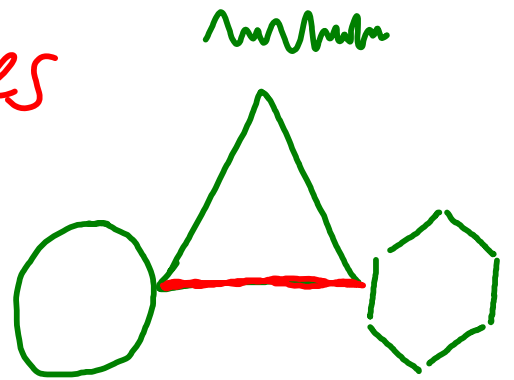


Thm. [Reverse AEP]:

If  $(Z_1, \dots, Z_n) \sim \text{Unif}(\text{Ball}(\underline{0}, \sqrt{n}\sigma^2))$ ,

then  $Z_1 \xrightarrow{\text{dist}} N(0, \sigma^2)$  as  $n \rightarrow \infty$

# Iso-perimetric Inequalities (Sphere bounds)



Ball minimizes  
\* \* \*  
over all bodies  
of a fixed volume!

$$\sigma^2(\Omega) \geq \sigma^2(\text{ball with radius } r_{\text{eff}})$$

$$P_e(\Omega) \geq P_e(\text{ " " " " })$$



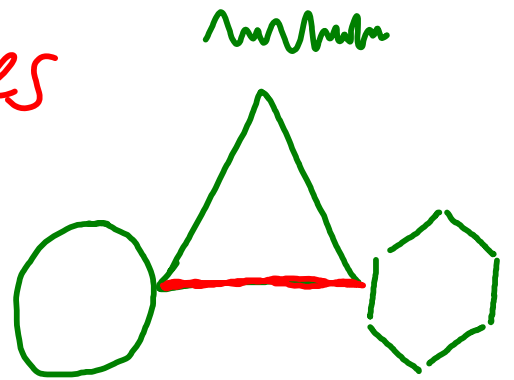
$$G(\Omega) \geq \text{N.S.M. of } n\text{-dim ball}$$

$$\mu(\Omega, P_e) \geq \text{V.N.R. " " " "}$$

# Iso-perimetric Inequalities

$$G(\Omega) \geq G_n(\text{Ball})$$

$$\mu(\Omega, \rho_e) \geq \mu_n(\text{Ball}, \rho_e)$$



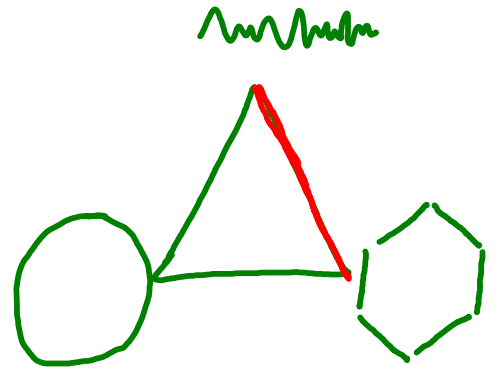
Asymptotic sphere bounds:

$$G_n(\text{Ball}) \rightarrow \frac{1}{2\pi e} \quad \text{as } n \rightarrow \infty$$

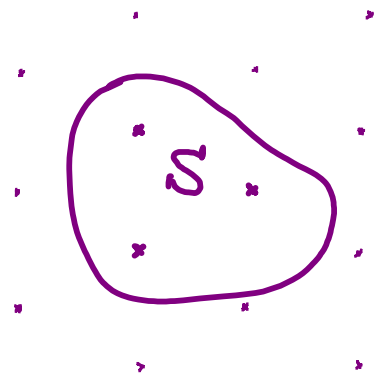
$$\mu_n(\text{Ball}, \rho_e) \rightarrow 2\pi e \quad \text{as } n \rightarrow \infty$$



# A Random Lattice Ensemble: Minkowski - Hlawka - Siegel



$N_{\Lambda}(S) \triangleq$  number of nonzero points of  $\Lambda$  inside a body  $S$



Theorem: For every dimension  $n$ , there exists an ensemble  $\{\Lambda\}$  of lattices with a constant point density  $\gamma = \frac{1}{V_{\Lambda}}$  (= a prob. measure over all generator matrices  $G$  with determinant  $1/\gamma$ ) such that for every bounded body  $S$

$$E_{MHS} \{ N_{\Lambda}(S) \} = \gamma \cdot \text{Vol}(S)$$

Just like a uniformly distributed random code!  $\nabla$

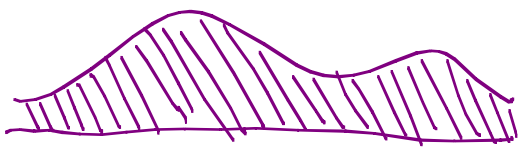
# Minkowski - Hlawka - Siegel



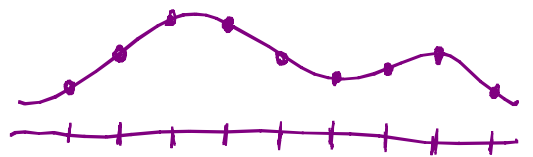
1. For any Riemann integrable function  $f(\cdot)$

$$\text{integral} = \frac{1}{\sqrt{v}} \cdot E_{\text{MHS}} \left\{ \begin{array}{c} \text{lattice-samples} \\ \text{sum} \end{array} \right\}$$

$$\int_{\mathbb{R}^n} f(x) dx$$



$$\sum_{\lambda \in \Lambda} f(\lambda)$$

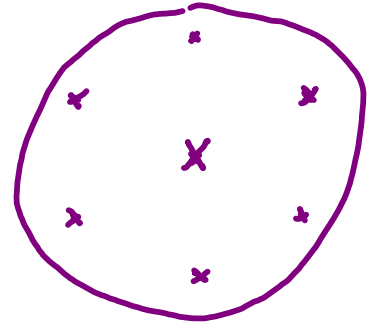


2. There exists (at least one) lattice which is (at least) as "good" as (1.)

# Implication 1 : packing Goodness

$$S = \text{Ball}(0, r)$$

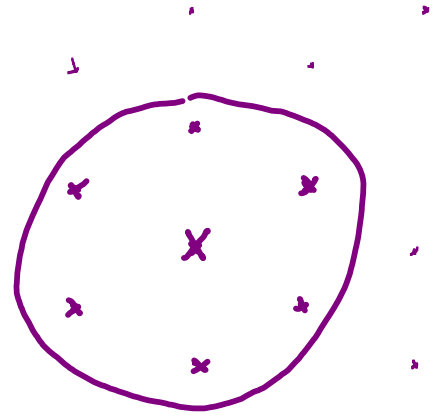
$$E_{\text{MHS}} \{N_{\mathcal{L}}(\text{Ball})\} = \gamma \cdot V_n \cdot r^n$$



# Implication 1 : packing Goodness

$$S = \text{Ball}(0, r)$$

$$E_{\text{MHS}} \{N_{\mathcal{L}}(\text{Ball})\} = \nu \cdot V_n \cdot r^n$$



$$\text{If } \text{Vol}(\text{Ball}) = V_n \cdot r^n < 1/\nu$$

$$\iff r < r_{\text{eff}} \quad (**)$$

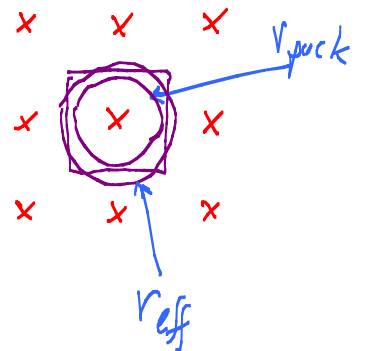
$$\implies E\{N_{\mathcal{L}}\} < 1$$

But  $N_{\mathcal{L}} = \text{integer}$

$$\implies N_{\mathcal{L}} = 0 \text{ for some } \mathcal{L} \in \text{MHS}$$

$$\implies d_{\min} = \|\text{shortest vector}\| > r$$

$$\implies r_{\text{pack}} > r/2 \quad (***)$$



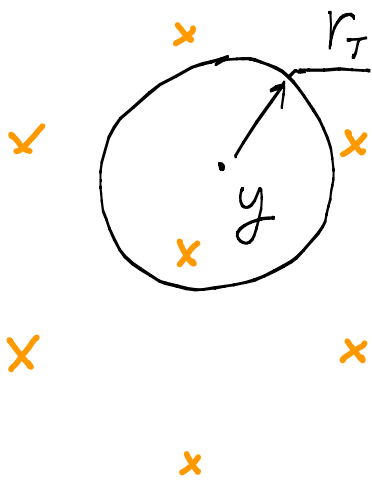
$$(**) + (***) \implies \text{packing efficiency} = \frac{r_{\text{pack}}}{r_{\text{eff}}} \geq 1/2$$

(for each dim  $n$ )

# Implication 2: Modulation goodness

Theorem (achieving Poltyrev capacity):

$$\exists \Lambda_n \text{ s.t. } \mu(\Lambda_n, P_e) \rightarrow 2\pi e, \quad n \rightarrow \infty.$$



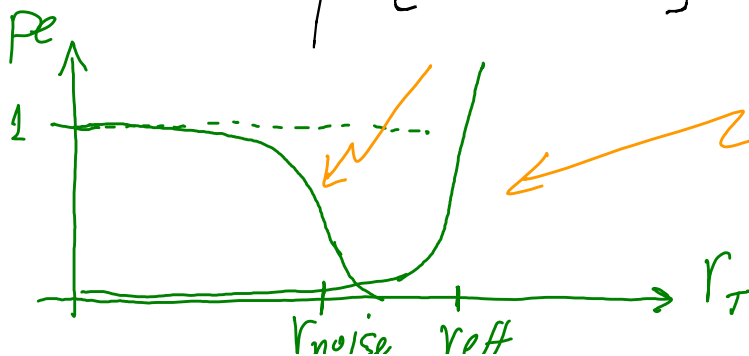
$$y = \lambda + z$$

$$\hat{\lambda} = \begin{cases} \Lambda \cap \text{Ball}(y, r_T), & \text{if unique} \\ ? , & \text{if empty or non-unique} \end{cases}$$

where  $r_T$  = search radius ("threshold")

$$\Rightarrow P_e = \Pr \left\{ \text{true } \lambda \text{ outside Ball OR other } \lambda \text{ inside Ball} \right\}$$

$$= \Pr \{ \|z\| > r_T \} + N_{\Lambda}(\text{Ball}(z, r_T))$$



$$\propto \frac{1}{V_n} \cdot r_T^n$$

$$\left( \frac{r_T}{r_{\text{eff}}} \right)^n$$

# Implication 3: Asymptotic Goodness

~ for Covering & Quantization ~

$\exists \Lambda_n$ :

$$f_{\text{cov}}(\Lambda_n) \triangleq \frac{V_{\text{cov}}(\Lambda_n)}{V_{\text{eff}}} \xrightarrow{n \rightarrow \infty} 1 \quad [\text{Rojers 1950}]$$

$$\Rightarrow G(\Lambda_n) \rightarrow \frac{1}{2\pi e} \quad [\text{Poltgyrev-Z-Feder 1995}]$$

Note:

1.  $\frac{1}{2\pi e} = \text{Normalized Second Moment (Ball)}$   
 $n \rightarrow \infty$

2.  $\frac{1}{2} \log(2\pi e G(\Lambda)) = D(\text{Unif}(V_0) \parallel \text{WGN})$

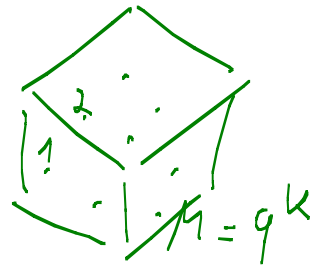
$\Rightarrow \text{cell} \rightarrow \text{Ball}, \text{unif}(\text{cell}) \rightarrow \text{WGN}$

# Alternative Ensemble:

Random Construction A (Loeliger 97, Erez et al 2005)

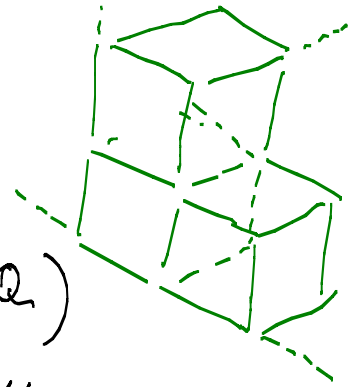
Let  $\mathcal{C} = q$ -ary  $(n, k)$  linear code over  $\mathbb{Q} = \{0, \dots, q-1\}$

$$= \{ \underset{n \times k}{G} \cdot \underline{i} : \underline{i} \in \mathbb{Q}^k \}$$



Let  $\Lambda_{\mathcal{C}} =$  modulo- $q$  lattice

$$= \{ \lambda \in \mathbb{R}^n : \lambda \bmod q \in \mathcal{C} \}$$



$G$  random (iid uniform on  $\mathbb{Q}$ )

$\Rightarrow \Lambda_{\mathcal{C}} =$  random lattice

$$\therefore G(\Lambda_{\mathcal{C}}), \mu(\Lambda_{\mathcal{C}}, P_e) = \text{func}\{q, k, n\}$$

# Simultaneous Goodness

Thm. [Erez-Litsyn-Z 2004]

There exists a sequence of lattices  $\Lambda_n$  in dim.  $n = 1, 2, \dots$ , such that as  $n \rightarrow \infty$

$$f_{\text{cov}}(\Lambda_n) \rightarrow 1$$

$$\underline{\lim} f_{\text{pack}}(\Lambda_n) \geq \frac{1}{2}$$

$$G(\Lambda_n) \rightarrow \frac{1}{2\pi e}$$

$$\mu(\Lambda_n, p_e) \rightarrow 2\pi e \quad \forall p_e > 0$$



# Random Construction A:

How Large  $q, k, n$  Should Be?

good packing:  $q, n \rightarrow \infty, 1 \leq k \leq n-1$   
 $f_{\text{pack}} \approx \frac{1}{2}$

good covering:  $q, k, n \rightarrow \infty$

$$f_{\text{cov}} \approx 1$$

good modulation:  
 $\mu(\mathcal{L}, p_e) \approx 2\pi e$

complete lattice

$$q, n \rightarrow \infty, 1 \leq k \leq n-1$$

finite constellation  
(= cosets)

$$n \rightarrow \infty, k \cdot \log(q) \ll n$$

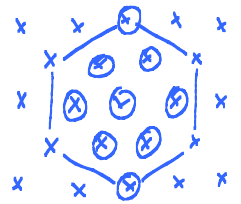
good quantization:  
 $G(\mathcal{L}) \approx \frac{1}{2\pi e}$

$$n \rightarrow \infty,$$

$$k \cdot \log(q) \ll n$$

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7. **Error exponents**
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11. Gaussian networks  $\text{Modulo}^n(\Lambda)$



# Tutorial-Part A Outline

## 7. Error exponents

$$E(\Lambda)$$

## Implication 2: Asymptotic Goodness for Modulation

$$P_e^{ML}(\mathcal{L}) = \int_0^{\infty} p(\|z\|=r) \cdot p(\text{nonzero codeword in Ball}(z, r)) dr$$

[Gallager 1962]

$\uparrow$   
 $N_{\mathcal{L}}(\text{Ball}(z, r))$

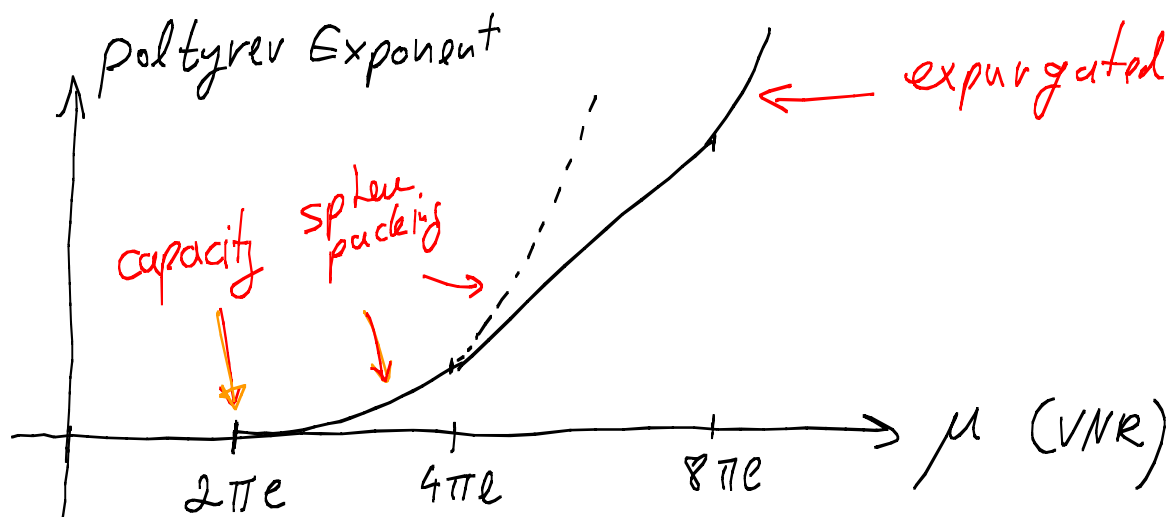
# Implication 2: Asymptotic Goodness for Modulation

$$P_e^{ML}(\mathcal{L}) = \int_0^{\infty} p(\|z\|=r) \cdot P(\text{nonzero codeword in Ball}(z, r)) dr$$

[Gallager 1962]

$$\approx N_{\mathcal{L}}(\text{Ball}(z, r))$$

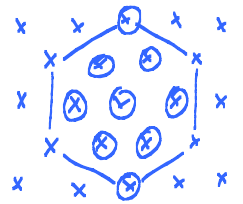
$$E_{MHS} \{ \} \Rightarrow r \cdot V_n \cdot r^n$$



$$\therefore \exists \mathcal{L}_n : \mu(\mathcal{L}_n, p_e) \xrightarrow{n \rightarrow \infty} 2\pi\epsilon \quad \forall p_e > 0$$

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# Tutorial-Part A Outline

## 8. Nested lattices

$$\Lambda_2 \subset \Lambda_1$$

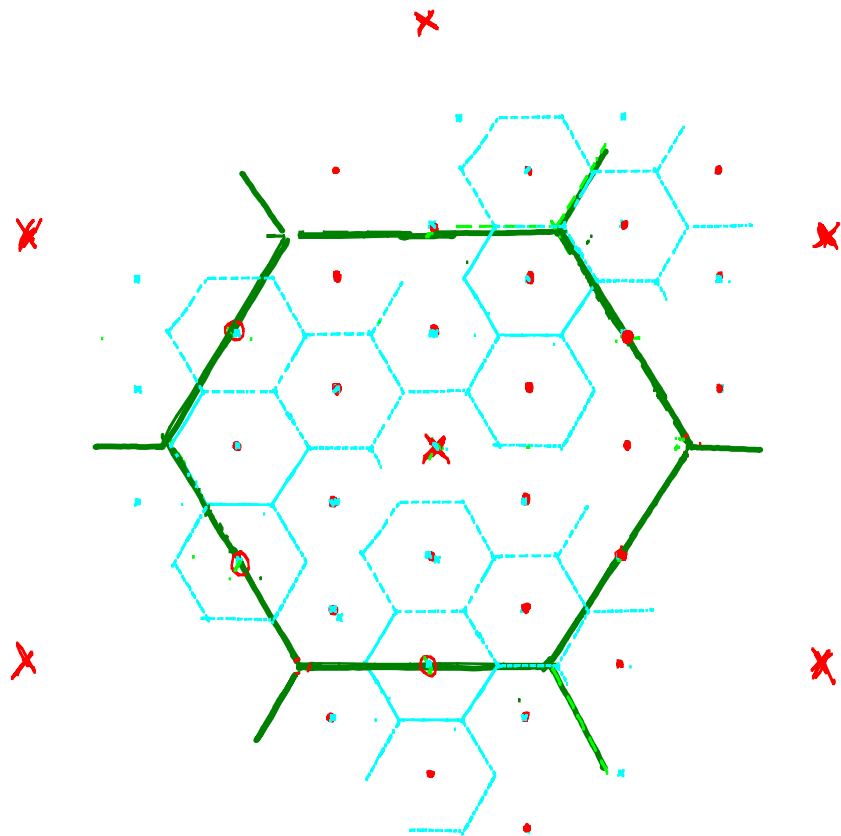
# Nested Lattices

$$\Lambda_2 \subset \Lambda_1 \Rightarrow \underline{G}_2 = \underline{G}_1 \cdot \underline{J}$$

course lattice      fine lattice      integer matrix

$$\text{Nesting Ratio} = \left( \frac{V_2}{V_1} \right)^{1/k} = |\det(\underline{J})|^{1/k}$$

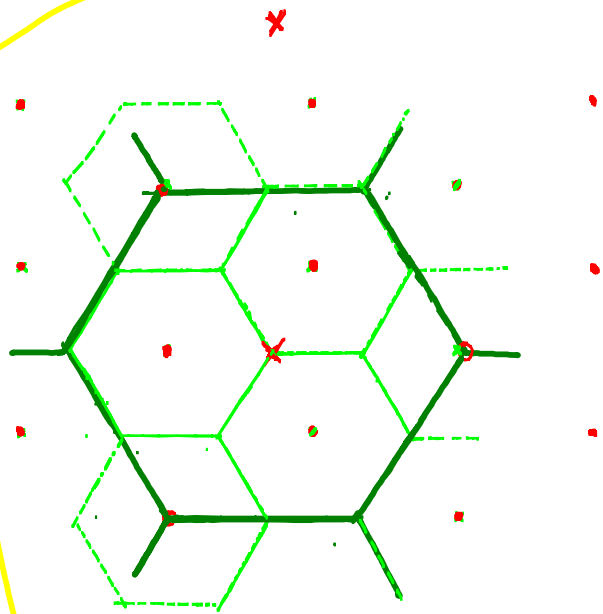
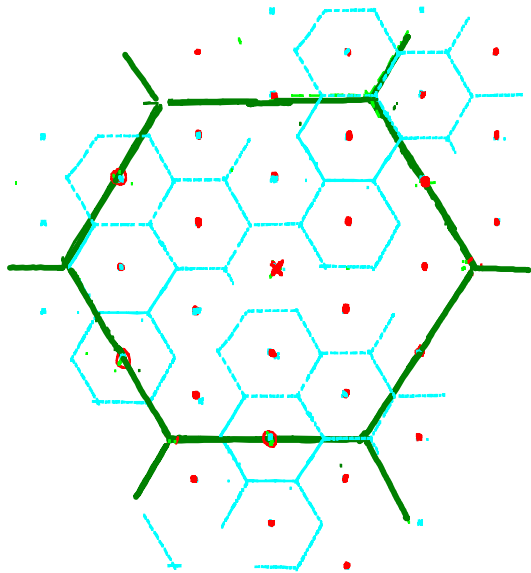
4:1





Not necessarily "Self Similar"!  
 $\Rightarrow V_2 \not\subset V_1$

Nested & Self Similar



Relatively periodic  
 (non nested)

## Diagonal Form

If  $\Lambda_2 \subset \Lambda_1$ , then  $\exists$  generator matrices  $G_1, G_2$   
 s.t. the nesting matrix  $J$  is diagonal

$$J = \begin{pmatrix} j_1 & & 0 \\ & \ddots & \\ 0 & & j_n \end{pmatrix}$$

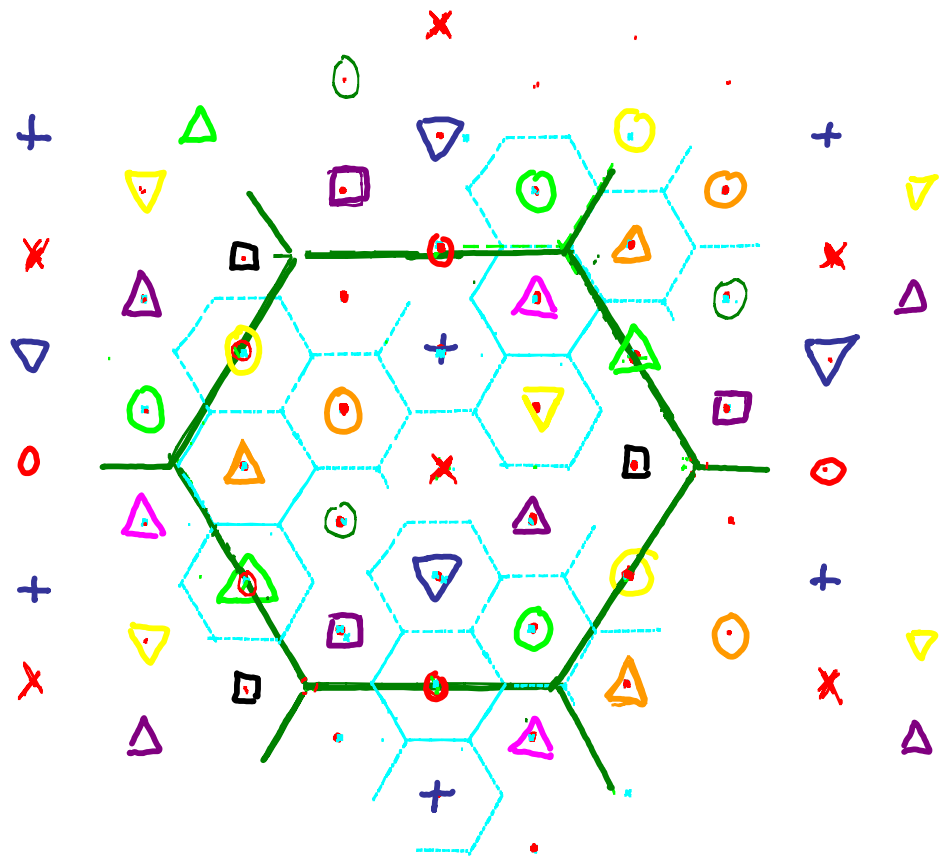
# Nested Lattices

$$\Lambda_2 \subset \Lambda_1 \Rightarrow \underline{\underline{G_2}} = \underline{\underline{G_1}} \cdot \underline{\underline{J}}$$

Relative Cosets =  $\Lambda_2 / \Lambda_1$

Coset  $\triangleq l_1 + \Lambda_2$ , for some  $l_1 \in \Lambda_1$

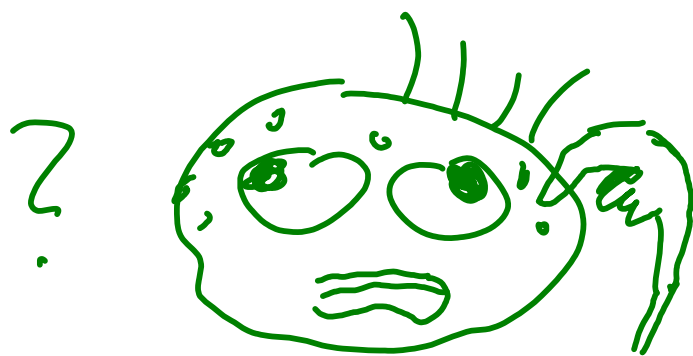
$$|\Lambda_2 / \Lambda_1| = V_2 / V_1 = |\det(J)|$$



Open Question :

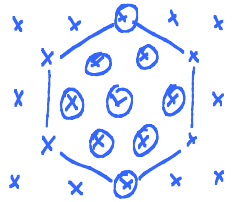
Good nested pairs ?

( low shaping & coding losses  $\Rightarrow$   
low  $G(\Lambda_2) * \mu(\Lambda_1, p_e)$  )



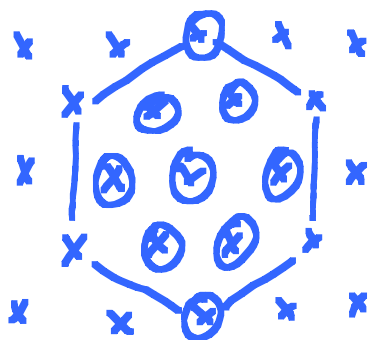
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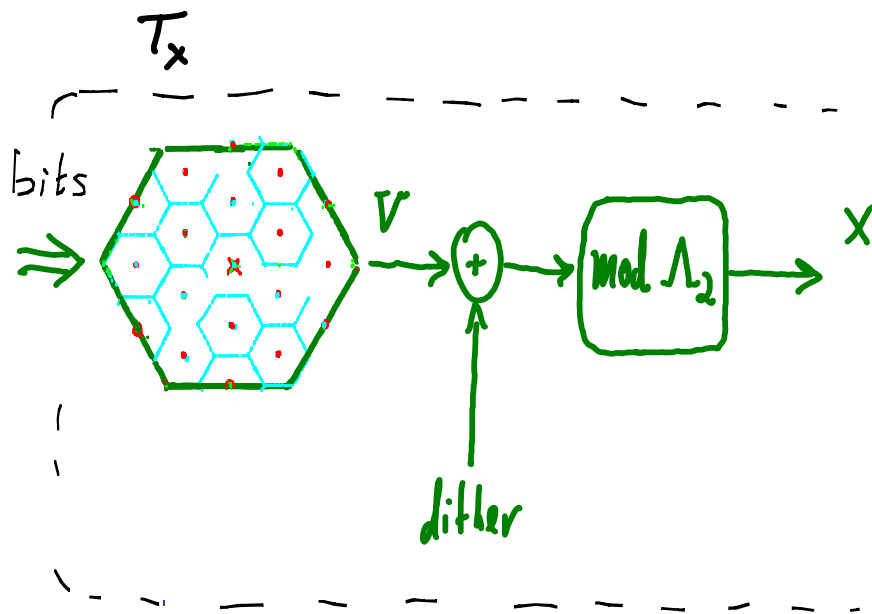
# Tutorial-Part A Outline

## 9. Lattice (Voronoi-) shaping



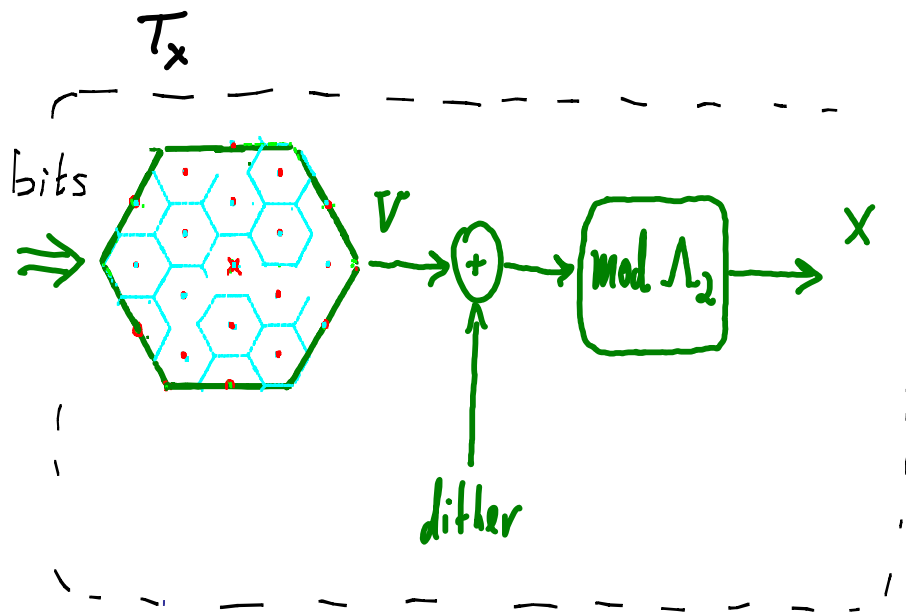
# Dithered Voronoi Modulation

message = fine point in coarse fund. cell  
 $\Leftrightarrow$  coset in  $\Lambda_1/\Lambda_2$



transmitted vector = reduction to Voronoi cell  
 $\Leftrightarrow$  min energy coset point

# Dithered Voronoi Modulation

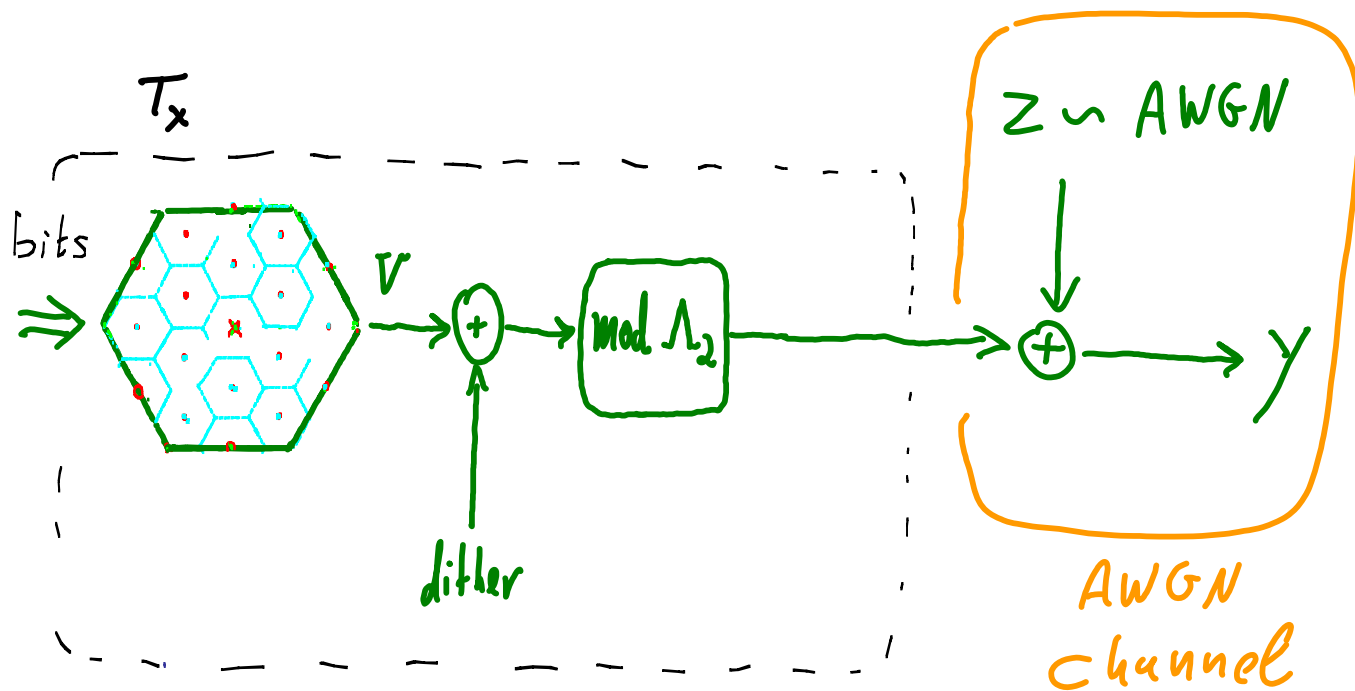


**Crypto Lemma**  $\Rightarrow$  for any  $v \in \mathcal{V}_\Lambda$ :

$$X = [v + \text{dither}] \bmod \Lambda_2 \sim \text{Unif}(\mathcal{V}_0)$$

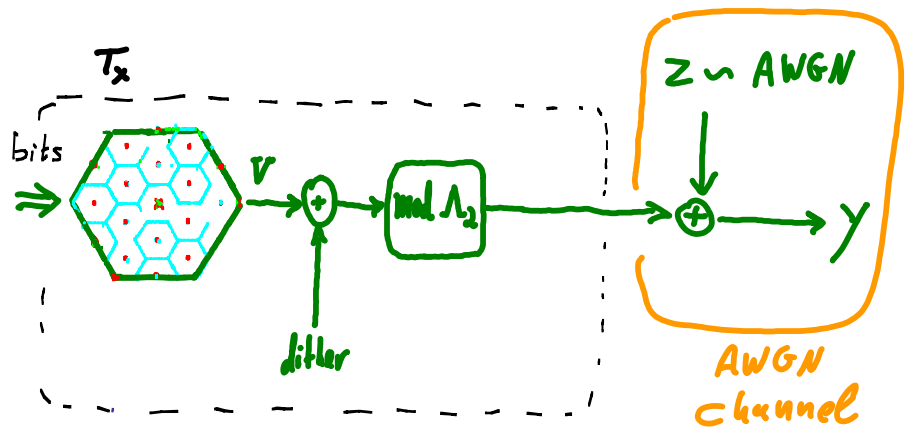
$$\Rightarrow T_x \text{ power} = \frac{1}{n} E \|X\|^2 = \sigma_{\Lambda_2}^2$$

# Dithered Voronoi Modulation: Decoding





# Dithered Voronoi Modulation: Decoding



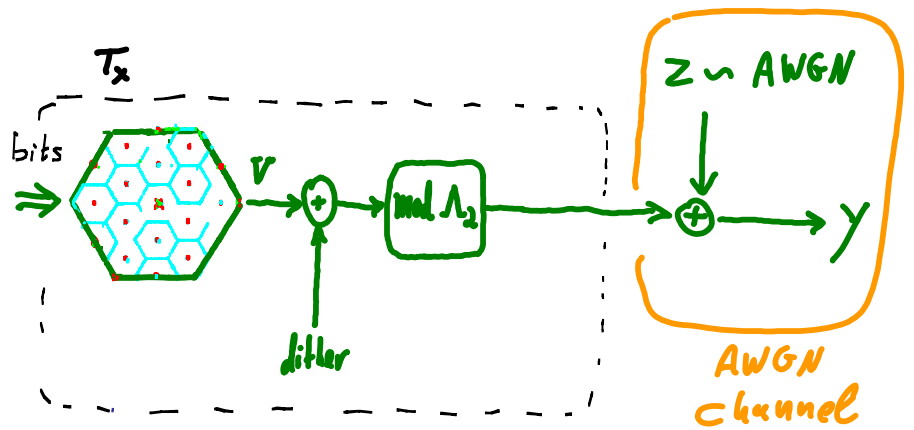
## Maximum-Likelihood Decoding:

$$\hat{v}_{ML} = \arg \max_{v \in \mathcal{V}_L} \|y - [v + \text{dither}] \bmod \Lambda_2\|$$

$\Rightarrow$  constrained search  $\Rightarrow$  complex !

reduction  
to %

# Dithered Voronoi Modulation: Decoding



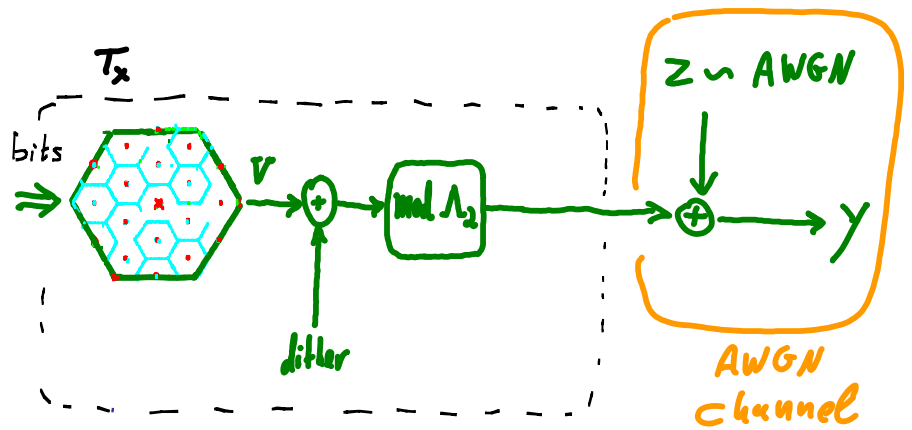
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$\Rightarrow$  constrained search  $\Rightarrow$  complex !

reduction  
to %

# Dithered Voronoi Modulation: Decoding



## Maximum-Likelihood Decoding:

$$\hat{\sigma}_{ML} = \arg \min_{\sigma \in \mathcal{L}_2} \|y - [\sigma + \text{dither}] \bmod \Lambda_2\|$$

⇒ constrained search ⇒ complex !

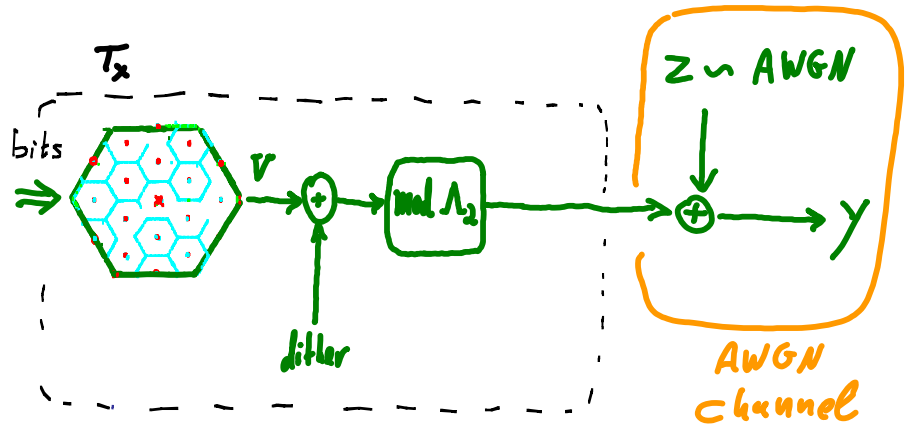
## Lattice Decoding:

$$\hat{\sigma}_{LD} = \left( \arg \min_{\lambda \in \Lambda_1} \|y - [\lambda + \text{dither}]\| \right) \bmod \Lambda_2$$

⇒ un-constrained search ⇒ easier ?

reduction  
to set of  
representatives

# Dithered Voronoi Modulation: Decoding



## Coset Decoding:

$$\hat{v}_{CD} = \underset{\text{coset} \in \Lambda_1/\Lambda_2}{\text{argmax}} \sum_{\lambda \in \text{coset}} e^{-\frac{\| \cdot \|^2}{2\sigma^2}}$$

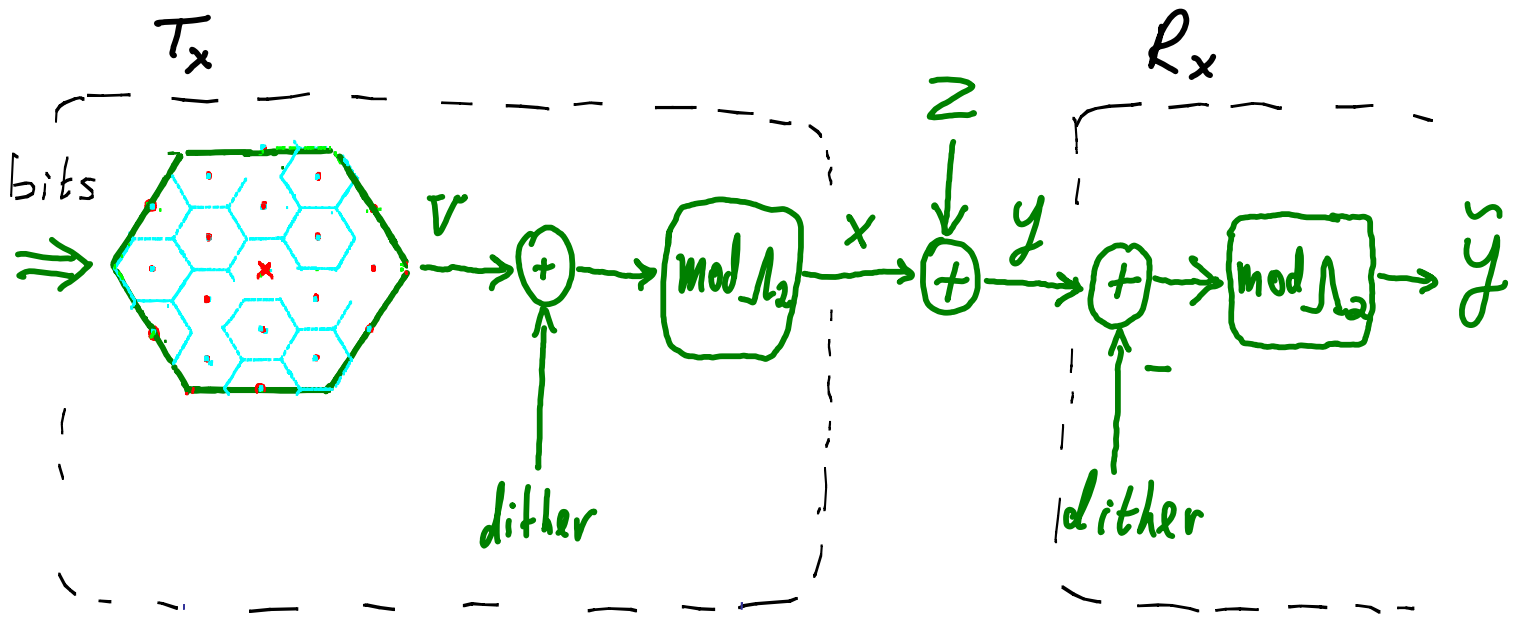
# Open Question 5:

Simple Decoding of  
Shaped Lattices ?

Remark : keep symmetry of  $p_{ei}$ , power: ?



# Dithered Voronoi Modulation: Equivalent Channel



$$T_x: X = [v + \text{dither}] \bmod \Lambda_2$$

$$\text{Channel: } Y = X + Z$$

$$R_x: \tilde{y} = [y - \text{dither}] \bmod \Lambda_2$$

$$\hat{v}_{LD} = Q_{\Lambda_1}(\tilde{y})$$

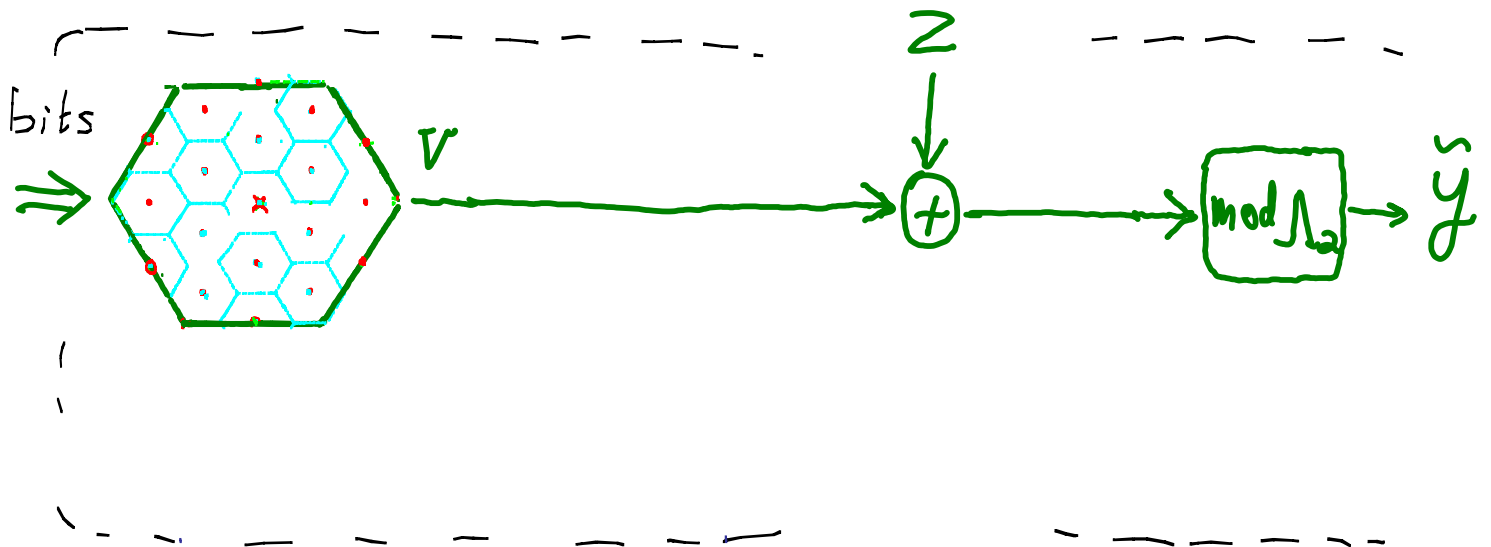
reduction  
to  $\mathcal{V}_0$

reduction  
to (any) set  
of coset rep.

# Dithered Voronoi Modulation: Equivalent Channel

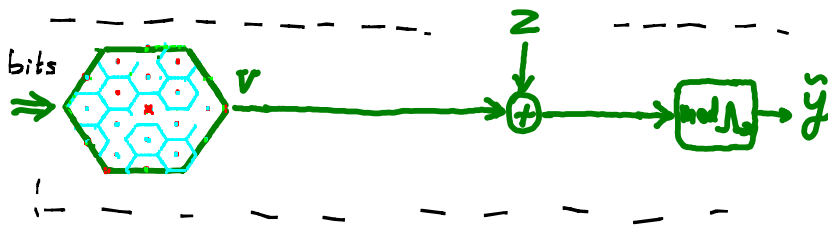
Distributive Property :

$$(a + b \bmod \Lambda) \bmod \Lambda = (a+b) \bmod \Lambda$$



# Dithered Voronoi Modulation : Performance

$\Lambda_1 =$  good for  $N(0, \sigma_z^2) \Rightarrow P_e < \epsilon \forall V$   
 $\Lambda_2 =$  good for quantization  $\Rightarrow \sigma_{\Lambda_2}^2$  small w.r.t  $V_2$



Rate =  $\frac{1}{n} \log \left( \frac{V_2}{V_1} \right)$  bit/channel use

NSM( $\Lambda_2$ )  
VNR( $\Lambda_1$ )

$= \frac{1}{2} \log \left( \frac{P}{\sigma_z^2} \right) - \frac{1}{2} \log \left( G(\Lambda_2) \cdot \mu(\Lambda_1, \epsilon) \right)$

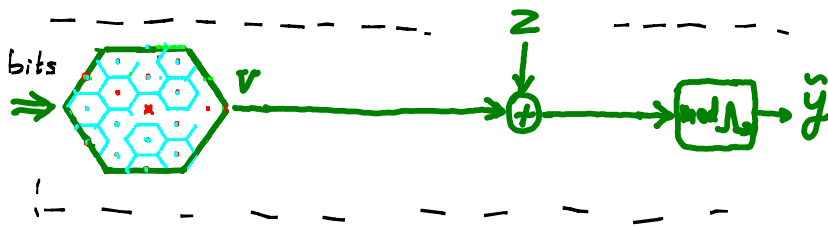
AWGN capacity @ High SNR

LOSS  $\rightarrow 0$   
 $n \rightarrow \infty$   
 for good lattices ....



# Dithered Voronoi Modulation : Performance

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AWGN capacity @ High SNR

$-\frac{1}{2} \log (G(\Lambda_2) \cdot \mu(\Lambda_1, \epsilon))$

LOSS  $\rightarrow 0$

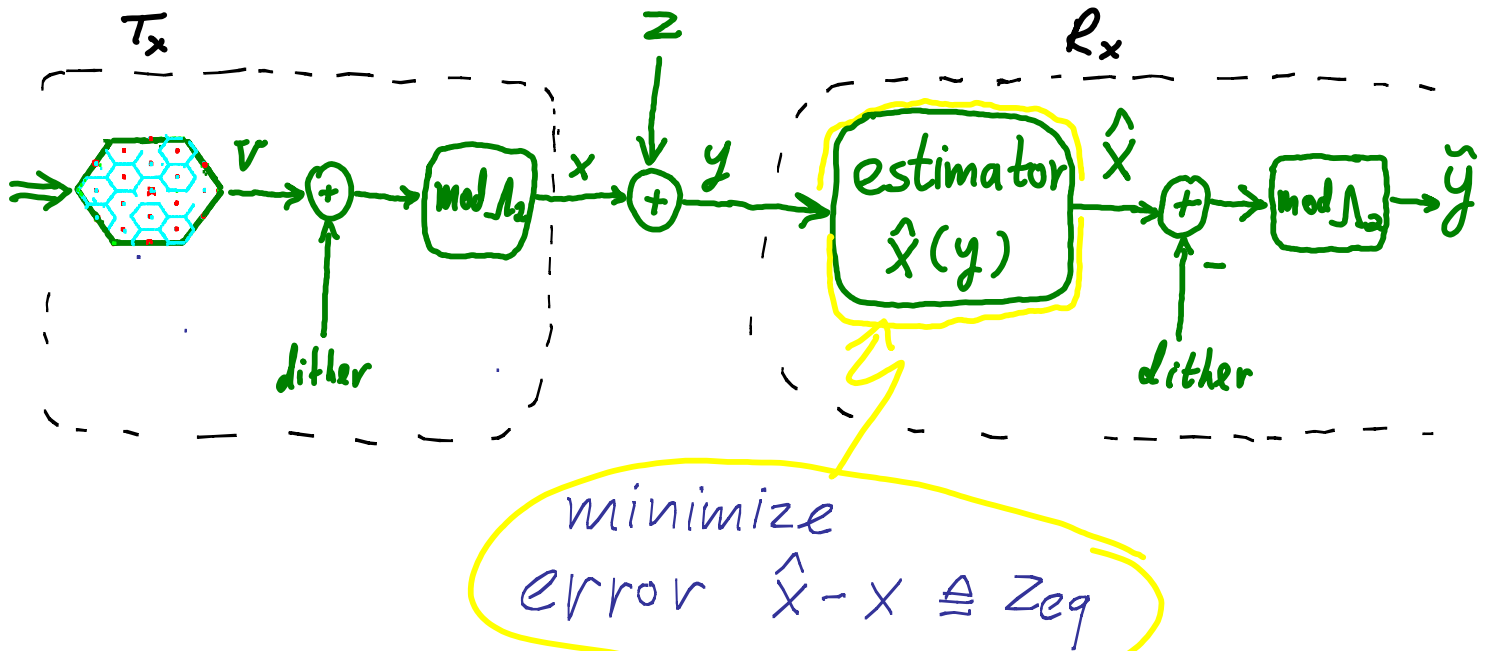
$n \rightarrow \infty$

for good lattices ....

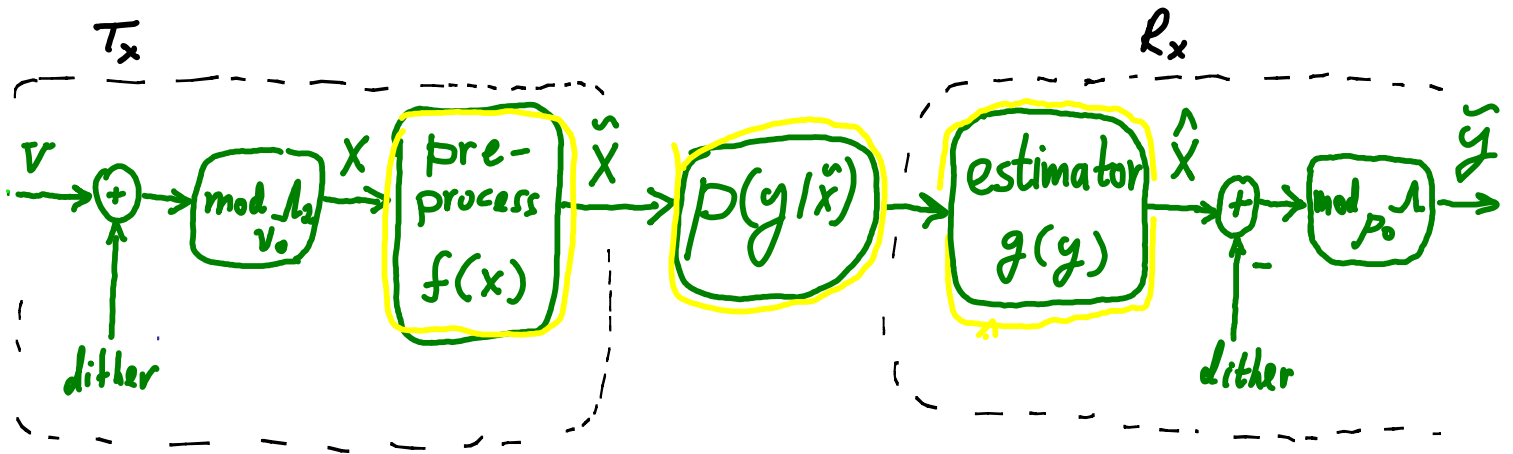
Shaping loss  $\triangleq \frac{1}{2} \log (2\pi e G(\Lambda_2))$

coding loss  $\triangleq \frac{1}{2} \log ( \mu(\Lambda_1, P_e) / 2\pi e )$

Achieving  $\frac{1}{2} \log(1 + \text{SNR})$  @ general SNR

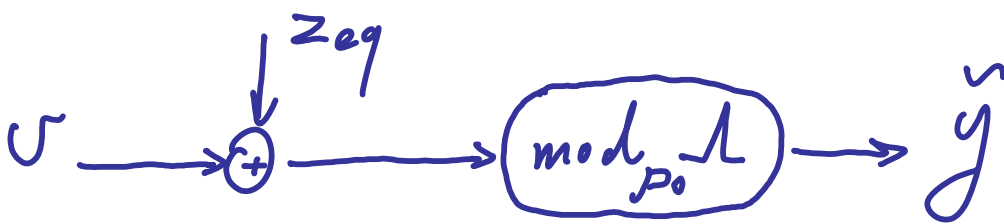


# The Equivalent Mod- $\Lambda$ Channel



## Thm. [Mod- $\Lambda$ channel]

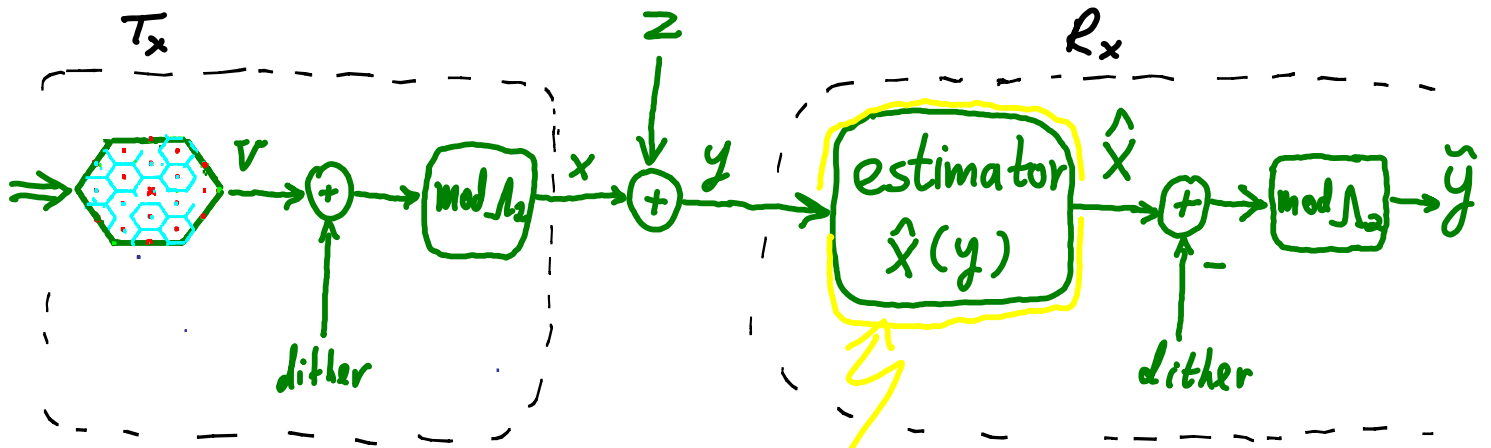
The channel  $v \rightarrow \tilde{y}$  is equivalent to a mod- $\Lambda$  channel



$$z_{eq} \stackrel{\text{dist.}}{=} \hat{X} - X \perp\!\!\!\perp v$$

$$X \sim \text{unif}(V_0), \quad \hat{X} = g(p(y|f(x)))$$

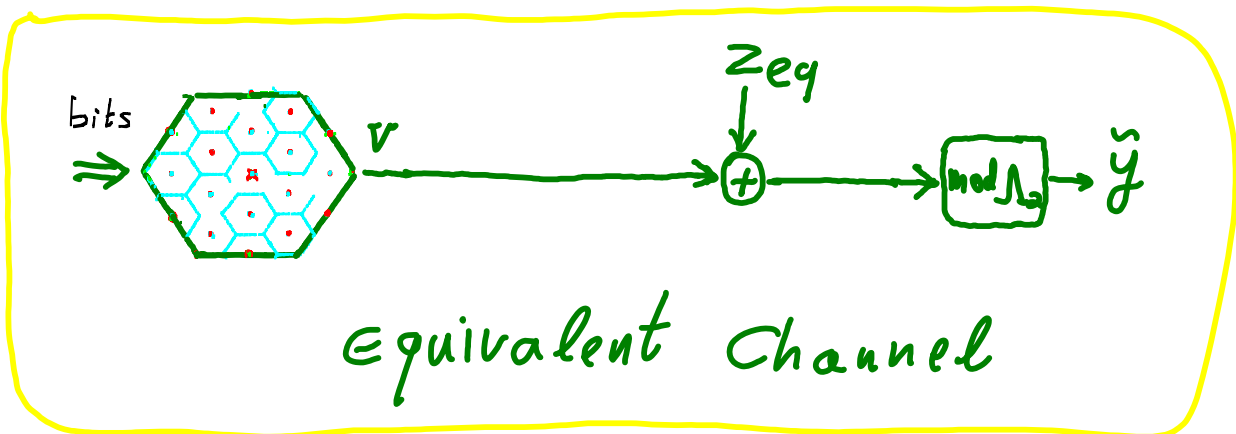
# Achieving $\frac{1}{2} \log(1 + \text{SNR})$ @ general SNR



minimize error  $\hat{x} - x \triangleq z_{eq}$

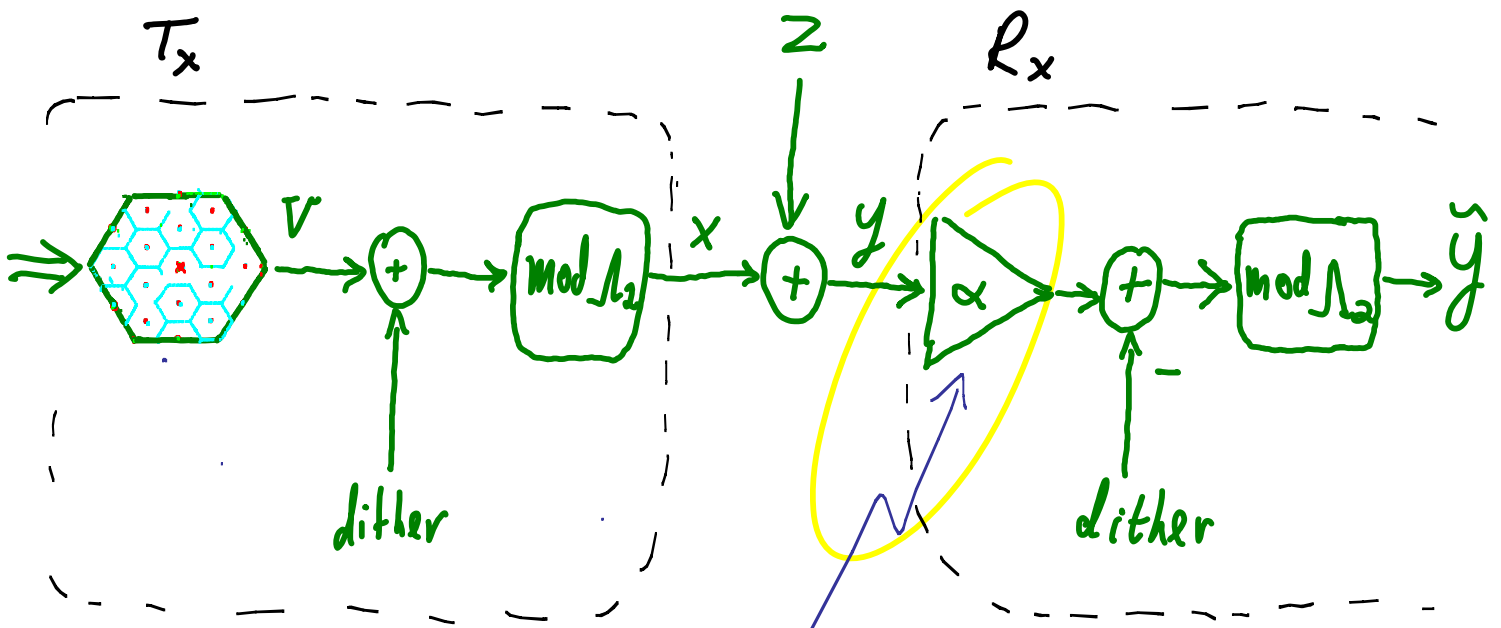
Crypto-Lemma  $\Rightarrow x \perp\!\!\!\perp v \Rightarrow z_{eq} \perp\!\!\!\perp v$

distributive property  $\Rightarrow \tilde{y} = (v + z_{eq}) \text{ mod } \Lambda_2$



Need to find  $\hat{x}(y)$  to optimize channel!

# Achieving $\frac{1}{2} \log(1 + \text{SNR})$ @ general SNR: Linear Estimation



$$\text{Var} \left\{ \underbrace{\alpha Z}_{\text{AWGN}} + \underbrace{(1-\alpha) \cdot \text{dither}}_{\text{"self noise"}} \right\} = \sigma_z^2 \cdot \frac{\text{SNR}}{1 + \text{SNR}}$$

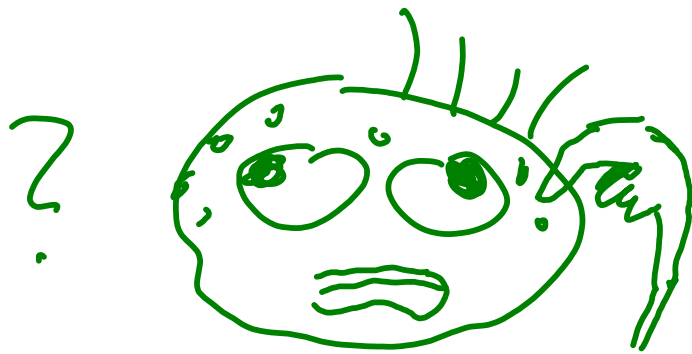
@  $\alpha = \alpha_{\text{wiener}}$

$$\Rightarrow \text{Rate} = \underbrace{\frac{1}{2} \log(1 + \text{SNR})}_{\text{AWGN capacity}} - \underbrace{\text{Shaping * Coding loss}}_{\substack{\rightarrow 0 \\ n \rightarrow \infty}}$$

# Open Question :

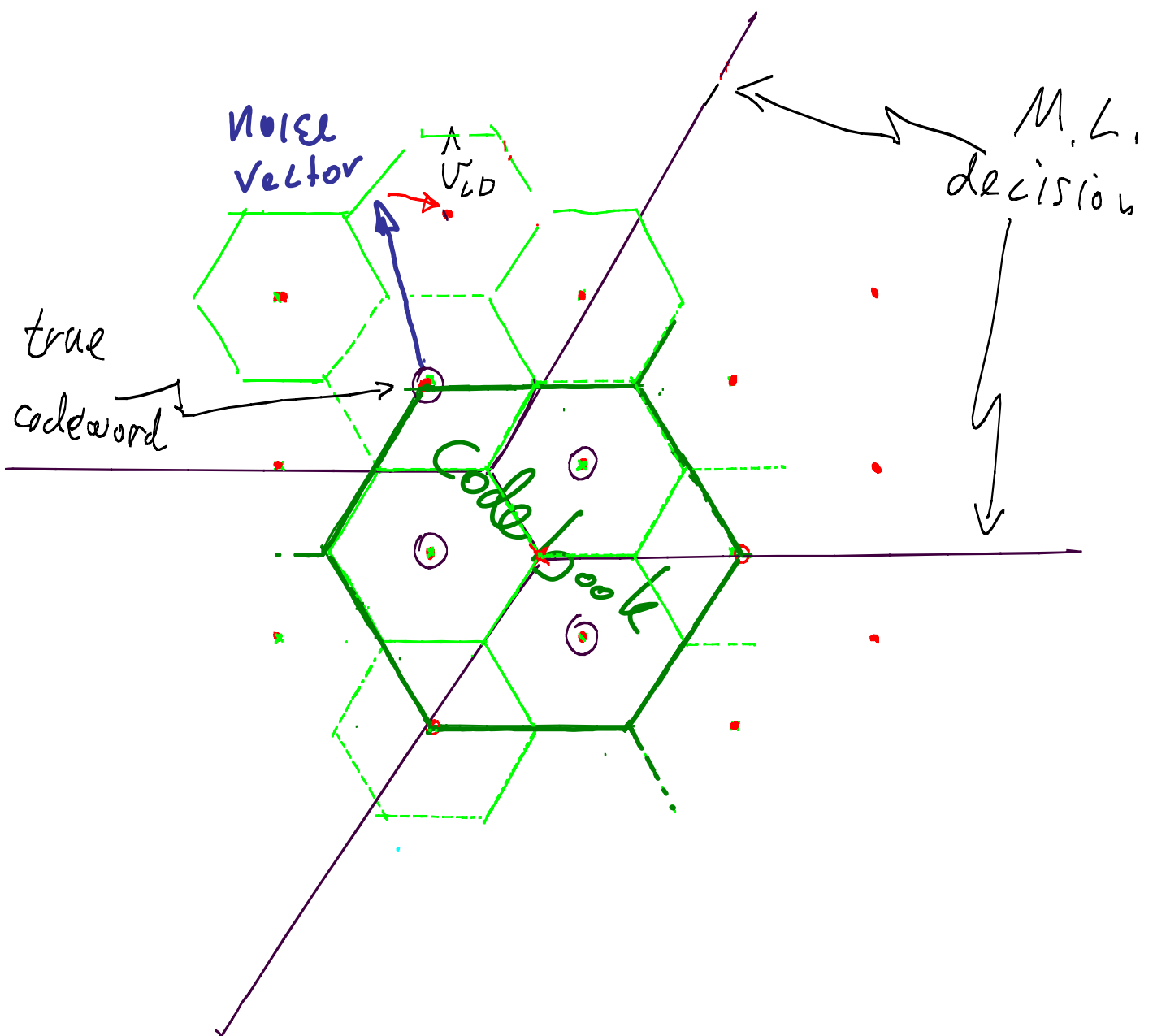
\* Bounds on **Volum to Noise Ratio** relative to mixture noise  $\mu(\Lambda, P_e, \alpha)$  ?

\* Good nested pairs @ general SNR ?

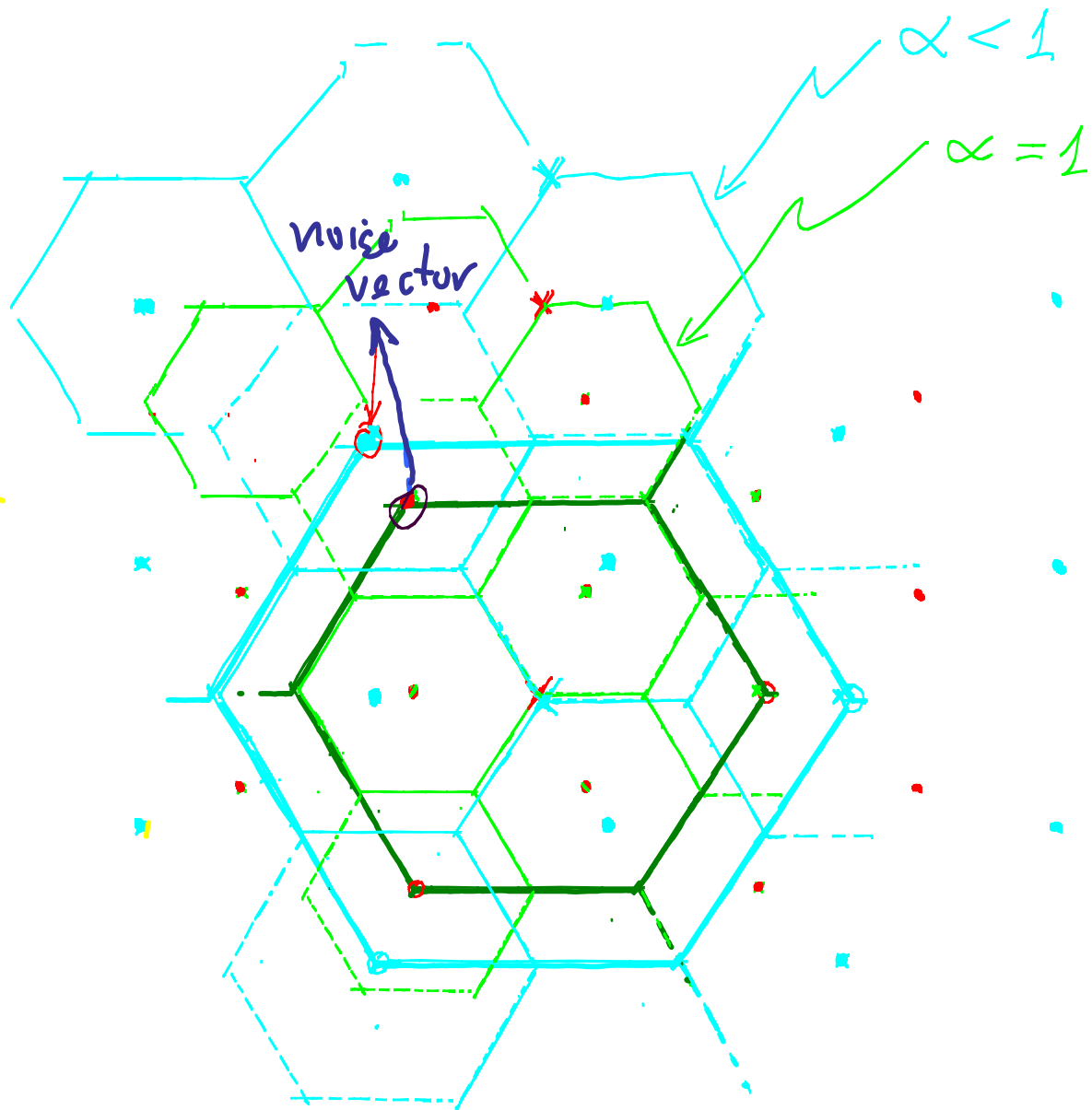


Geometric picture:

ML versus Lattice Decoding ( $\alpha=1$ )



Geometric picture :  
Linear Scaling ( $\alpha=1$ )  $\Rightarrow$  Lattice Inflation



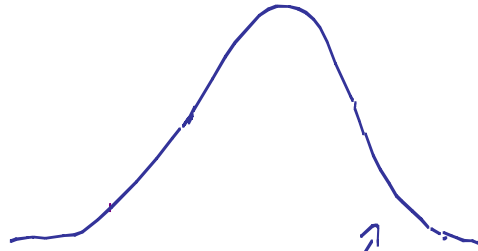
Lattice @ Decoder Inflated by  $\frac{1}{\alpha}$



# But Noise is Not Quite Gaussian ...

$$\alpha = 1$$

pure Gaussian →

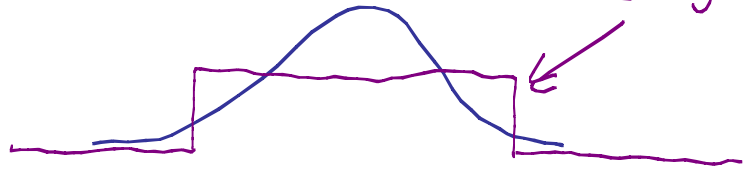


Gaussian noise

"self noise"

$$\alpha = \alpha_{\text{Wiener}}$$

minimum-energy  
mixture → →



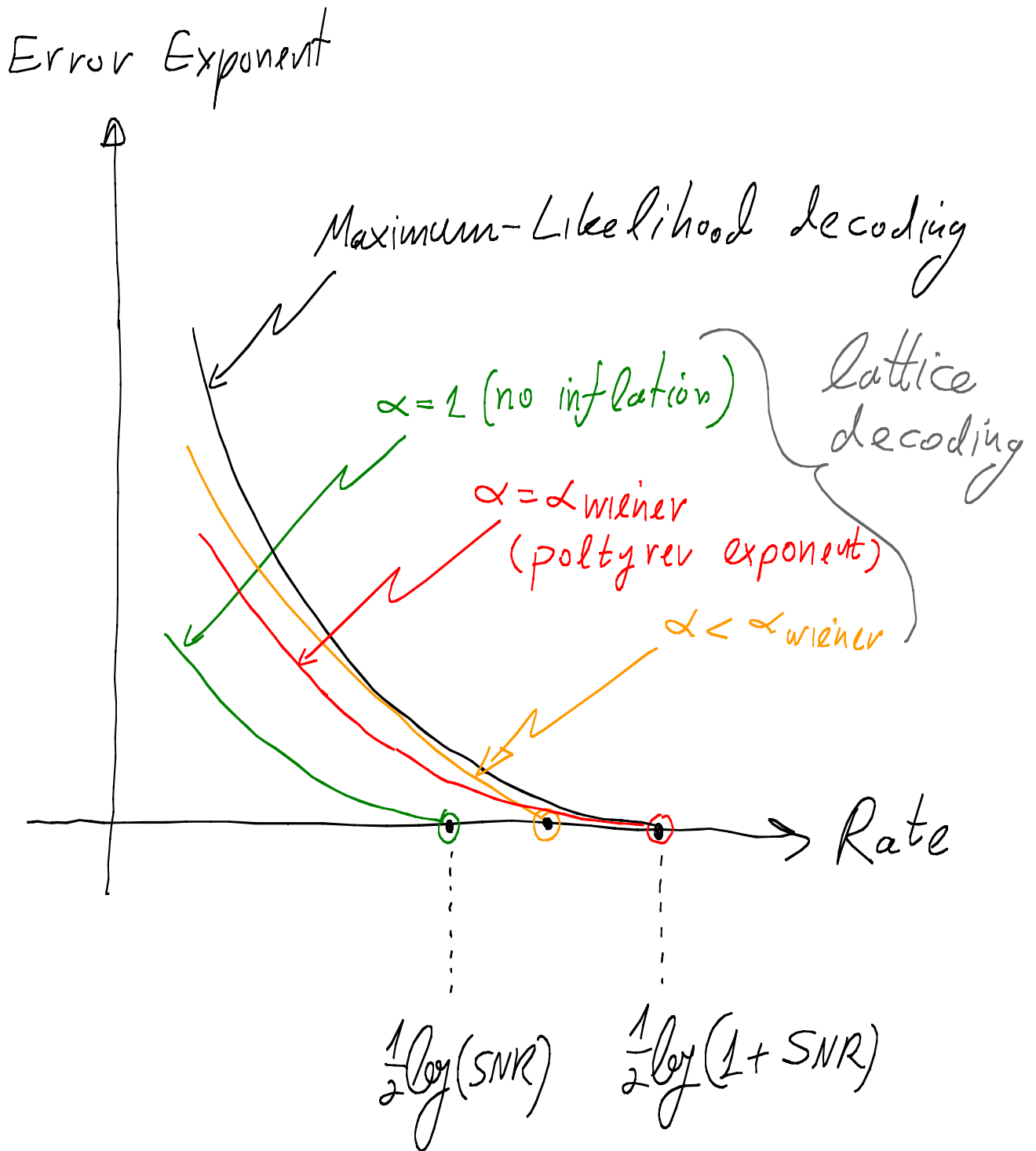
$$\alpha < \alpha_{\text{Wiener}}$$

reduced-tail  
mixture

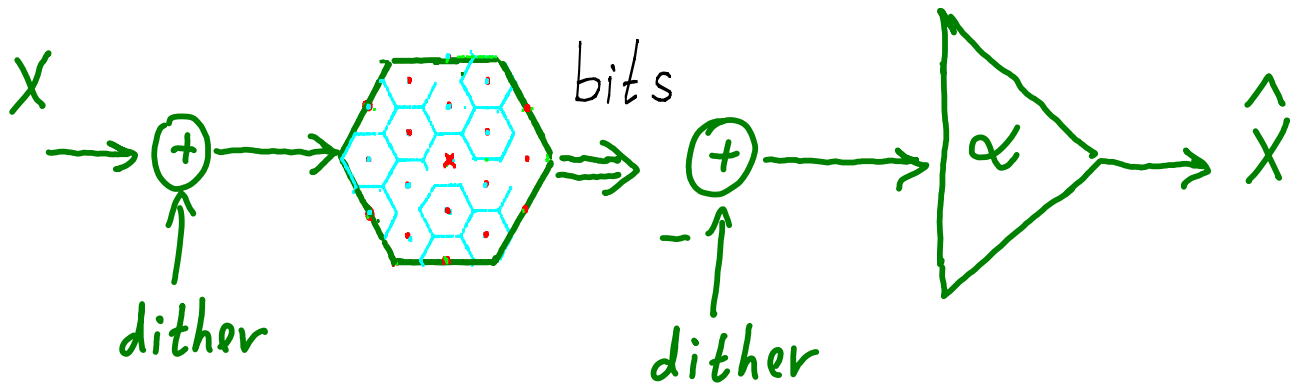


⇒ inflation-coefficient  $\alpha$  affects  
equivalent noise distribution

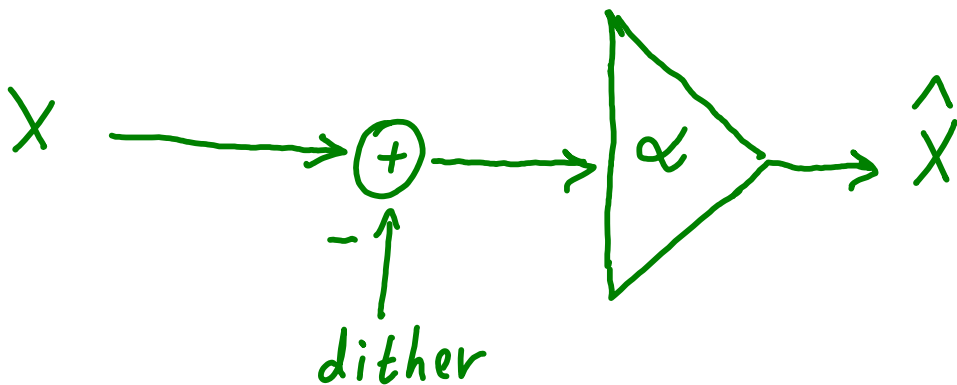
# Effect of $\alpha$ on Lattice-Decoding



# Voronoi Quantization



Crypto Lemma  $\Rightarrow$  Equivalent channel



$$\text{Distortion} = \text{Var} \{ \alpha \cdot \text{dither} + (1-\alpha)X \} = \frac{\sigma_x^2 \sigma_d^2}{\sigma_x^2 + \sigma_d^2}$$

@  $\alpha = \alpha_{\text{Wiener}}$

$$\text{Rate} = \underbrace{\frac{1}{2} \log \left( \frac{\sigma_x^2}{D} \right)}_{\text{Q-Gaussian RDF}} - \underbrace{\frac{1}{2} \log \left( G(\mathcal{L}_1) \cdot \mu(\mathcal{L}_2, P_e, \alpha) \right)}_{\text{loss} \rightarrow 0 \text{ as } n \rightarrow \infty}$$

Next file ... →

