# State-Dependent Channels with Composite State Information at the Encoder

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# Channel Model: Memoryless State-dependent Channel



Memoryless channel:

$$p(\mathbf{y}|\mathbf{x},\mathbf{s}) = \prod_{i=1}^{n} p(y_i|x_i,s_i).$$

"Memoryless" (i.i.d.) state sequence:

$$p(\mathbf{s}) = \prod_{i=1}^n p(s_i) \, .$$

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## State not Known



• Problem reduces to "regular" DMC:

$$p(y|x) = \sum_{s \in S} p(s)p(y|x, s)$$
$$p(\mathbf{y}|\mathbf{x}) = \prod_{i=1}^{n} p(y_i|x_i).$$

• Capacity [Shannon '48]:  $\max_{p_x} I(X; Y)$ 

## State Known at the Receiver



- State (S) available at Rx can be regarded as part of output.
- Define  $\tilde{Y} = (Y, S)$ .  $\downarrow$ Channel from X to  $\tilde{Y}$  with no state available at Rx.
- No special treatment is required for state available at Rx.

## State Known at the Transmitter



### Causal state knowledge [Shannon '58]

•  $x_i = \operatorname{func}(w, s_1^i) - \operatorname{at} \operatorname{time} i$ , only states  $s_1, \ldots, s_i$  are known.

Capacity: max<sub>p(t)</sub> I(T; Y) ,

where  $t : S \to X$ , i.e., mappings x = t(s).

• Equivalent representation:  $\max_{p(u),x(u,s)} I(U; Y)$ 

(U is independent of S)

## State Known at the Transmitter

Non-causal state knowledge [Gel'fand & Pinsker '80]

- $x_i = \operatorname{func}(w, s_1^n) \operatorname{at} \operatorname{time} i$ , all states  $s_1, \ldots, s_n$  are known.
- **Capacity:**  $\max_{p(u|s),x(u,s)} I(U;Y) I(U;S).$
- In terms of "Shannon strategies":  $x = t_s(s)$ .

(here random strategy t depends on s)

• Achievable using "random-binning".

NC knowledge with limited look-ahead [Weissman & El Gamal '06]

• Limited look-ahead k.

• 
$$x_i = \operatorname{func}\left(w, s_1^{i+k}\right)$$
 – at time *i*, states  $s_1, \ldots, s_{i+k}$  are known.

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## Composite State Known at the Transmitter

### Observations

- Achievables of causal and non-causal cases can be derived in the same way.
- Converses of causal and non-causal cases are similar and can be combined.

### State with parts known at Tx causally and non-causally

• 
$$S = (^{c}S, ^{nc}S).$$

- $^{c}S$  known *causally* at Tx.
- ${}^{nc}S$  known *non-causally* at Tx.
- Capacity:  $C = \max_{p(u|n^{c}s), x(u,n^{c}s,s)} [I(U;Y) I(U;n^{c}S)].$
- U is independent of <sup>c</sup>S given  ${}^{nc}S$ .

## Converse

$$n(R-\epsilon_n) \stackrel{\text{Fano}}{\leq} I(W;Y_1^n) = \sum_{i=1}^n I(W;Y_i|Y_1^{i-1}) \leq \sum_{i=1}^n I(W,Y_1^{i-1};Y_i)$$

## Shannon (causal)

• 
$$U_i \triangleq (W, Y_1^{i-1}).$$

- Causality ∜
  - $U_i$  is independent of  $S_i$ .

$$\sum I(W, Y_1^{i-1}; Y_i)$$

$$= \sum I(U_i; Y_i)$$

$$\leq n \max_{p(u), p(x|u,s)} I(U; Y)$$

er (non-causal)

• 
$$U_i \triangleq (W, Y_1^{i-1}, S_{i+1}^n)$$

- $S_i$  independent of  $(W, S_{i+1}^n)$
- Chain-rule for mutual informations.

$$\sum I(W, Y_1^{i-1}; Y_i) = \cdots =$$
$$= \sum I(U_i; Y_i) - I(U_i; S_i)$$

$$\leq \max_{p(u|s),p(x|u,s)} \left\{ I(U;Y) - I(U;S) \right\}$$

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## Converse

### In both converses...

- Achievable rate bounded by  $I(U_i; Y_i) I(U_i; S_i)$ .
- Auxiliary variable:  $U_i = (W, Y_1^{i-1}, S_{i+1}^n)$ .
  - Causal case: reduces to U<sub>i</sub> = (W, Y<sub>1</sub><sup>i-1</sup>), since S<sup>n</sup><sub>i+1</sub> is independent of Y<sub>i</sub>.

## • Can be used for the composite causal-non-causal case!

### Converse for the composite causal-non-causal case

- $S = ({}^{c}S, {}^{nc}S).$
- Achievable rate is bounded by  $I(U_i; Y_i) I(U_i; {}^{nc}S_i)$ .
- Auxiliary variable:  $U_i = (W, Y_1^{i-1}, {}^{\mathsf{nc}}S_{i+1}^n).$
- Maximization over  $p(u|^{nc}s)$  and  $x = \text{func}(u, {}^{c}s, {}^{nc}s)$ .

(*U* is independent of <sup>c</sup>S given  ${}^{nc}S$ )

# Achievable

## Causal (Shannon) case

- x = t(s).
- Strategy t is generated from W (input to equivalent channel).
- Can be thought of as "degenerated random binning" (U independent of S).

## Non-causal (Gel'fand-Pinsker) case

• Use random binning w.r.t. p(u|s) and x(u,s).

#### Composite causal-non-causal case

- Use random binning w.r.t  $p(u, {}^{nc}s)$  and  $x(u, {}^{c}s, {}^{nc}s)$ .
- Alternatively, combine random strategies w.r.t <sup>c</sup>S and random binning w.r.t. <sup>nc</sup>S.

Model SI Scenarios Composite SI@Tx WZ Compound

# Source Coding with Side-information at the Receiver



• X – Source.

- $\hat{X}$  Reconstructed (distorted) source at the decoder.
- S Side-information.
- $d(X, \hat{X})$  Distortion measure.

#### Side-information at the encoder ( $\mathbb{A}$ is closed)

- S can be regarded as part of the source  $(\tilde{X} \triangleq (X, S))$ .
- Distortion measure is w.r.t. X (and not  $\tilde{X}$ ).

# Source Coding with Side-information at the Receiver

Non-causal side-information [Wyner-Ziv '76]

$$R(D) = \min \left[ I(U; X) - I(U; S) \right],$$

where minimum is over all  $f : \mathcal{U} \times S \to \hat{\mathcal{X}}$ , s.t.  $E[d(X, f(U, S))] \leq D$ .

Causal side-information [Weissman-El Gamal '06]

 $R(D) = \min I(U; X)$ 

minimum over the same set as in the non-causal problem.

Composite causal-non-causal side-information

$$R(D) = \min \left[ I(U; X) - I(U; {}^{\mathsf{nc}}S) \right]$$

minimum over the same set.

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## When do Both Converses Diverge?

### Compound state-dependent channel:



•  $\theta \in \Theta$  – "Compound parameter": constant, unknown to Tx.

#### Compound channel with no state knowledge ( $\mathbb{A}$ is open)

- Worst-case capacity Maximal rate for all  $\theta$  simultaneously.
- Capacity: [Blackwell et al. '59; Dobrushin '59; Wolfowitz '60]

$$C^{\mathsf{wc}} = \max_{p(x)} \min_{\theta \in \Theta} I(X; Y)$$

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# When do Both Converses Diverge?

### Compound Channel with state S known *causally* at Tx

Trying to generalize Shannon's converse to compound case:

- $U_i \triangleq (W, Y_1^{i-1}).$
- U<sub>i</sub> depends on the statistics of Y which is unknown! (since θ not known).

$$n(R-\epsilon_n) \leq \sum_{i=1}^n I_{\theta}(W, Y_1^{i-1}; Y_i) \leq \sum_{i=1}^n I_{\theta}(W, Y_1^{i-1}; Y_i) \leq \sum_{i=1}^n I_{\theta}(W, Y_1^{i-1}; Y_i)$$

- U<sub>i</sub> ≜ (W, S<sub>1</sub><sup>i-1</sup>) No knowledge of θ is assumed! (Original U<sub>i</sub> used by Shannon)
- (Worst-case) Capacity: [Khina, Erez '10]

$$C^{\mathrm{wc}} = \max_{p(u), x(u,s)} \min_{\theta \in \Theta} I_{\theta}(U; Y)$$

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# When do Both Converses Diverge?

Compound Channel with state S known *non-causally* at Tx

• Similarly one would expect:

$$C^{\mathsf{wc}} = \max_{p(u|s), x(u,s)} \min_{\theta \in \Theta} I_{\theta}(U; Y)$$

- This is achievable but not optimal! [Piantanida, Shamai '10; Nair, El Gamal, Chia '10]
  - Marton's broadcast technique improves performance.

## Summary

- Different side-information scenarios can be treated similarly.
- Recognizing these similarities allows to:
  - Solve several different scenarios at once.
  - Combining results for composite/mixed scenarios.
- In more complex scenarios similar/combined treatments might diverge.