

Source Coding with Composite Side Information at the Decoder

Anatoly Khina

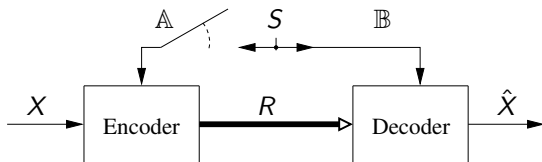
Joint work with: Uri Erez

Tel Aviv University

November 15, 2012

IEEEI, Eilat

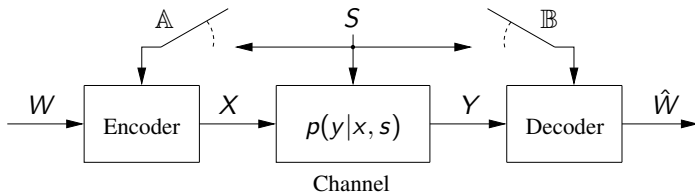
Source Coding with Side-information at the Decoder



- X – Source
- \hat{X} – Reconstructed (distorted) source at the decoder
- S – Side-information
- $d(X, \hat{X})$ – Distortion measure
- **Memorylessness (i.i.d.):**

$$p(\mathbf{x}, \mathbf{s}) = \prod_{i=1}^n p(x_i, s_i)$$

Dual Problem: Memoryless State-dependent Channel



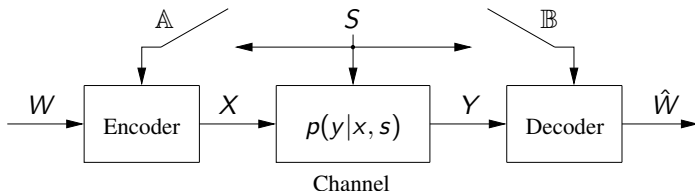
Memoryless channel:

$$p(\mathbf{y}|\mathbf{x}, \mathbf{s}) = \prod_{i=1}^n p(y_i|x_i, s_i).$$

“Memoryless” (i.i.d.) state sequence:

$$p(\mathbf{s}) = \prod_{i=1}^n p(s_i).$$

State not Known



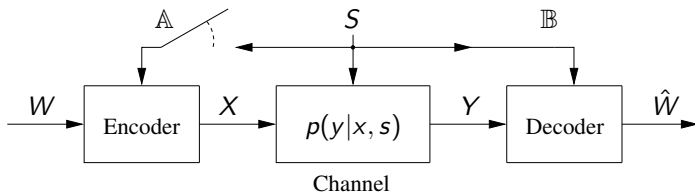
- Problem reduces to “regular” DMC:

$$p(y|x) = \sum_{s \in \mathcal{S}} p(s) p(y|x, s)$$

$$p(\mathbf{y}|\mathbf{x}) = \prod_{i=1}^n p(y_i|x_i).$$

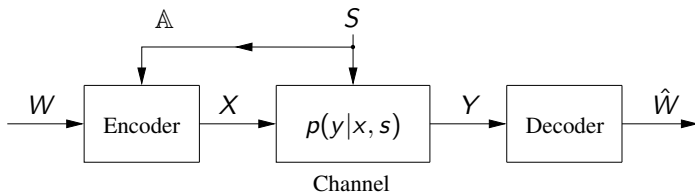
- **Capacity** [Shannon '48]: $\max_{p_x} I(X; Y)$

State Known at the Decoder



- State (S) available at Rx can be regarded as part of output
- Define $\tilde{Y} = (Y, S)$
 \Downarrow
 Channel from X to \tilde{Y} with no state available at Rx
- No special treatment is required for state available at Rx

State Known at the Encoder



Causal state knowledge [Shannon '58]

- $x_i = \text{func}(w, s_1^i)$ – at time i , only states s_1, \dots, s_i are known
- **Capacity:** $\max_{p(t)} I(T; Y)$,
where $t: \mathcal{S} \rightarrow \mathcal{X}$, i.e., mappings $x = t(s)$
- Equivalent representation: $\max_{p(u), x(u, s)} I(U; Y)$
(U is independent of S)

State Known at the Encoder

Non-causal state knowledge [Gel'fand & Pinsker '80]

- $x_i = \text{func}(w, s_1^n)$ – at time i , all states s_1, \dots, s_n are known
- **Capacity:** $\max_{p(u|s), x(u,s)} I(U; Y) - I(U; S)$
- In terms of “Shannon strategies”: $x = t_s(s)$
(here random strategy t depends on s)
- Achievable using “random-binning”

NC knowledge with limited look-ahead [Weissman & El Gamal '06]

- Limited look-ahead k
- $x_i = \text{func}(w, s_1^{i+k})$ – at time i , states s_1, \dots, s_{i+k} are known

Composite State Known at the Encoder

Observations

- Achievables of causal and non-causal cases can be derived in the same way
- Converses of causal and non-causal cases are similar and can be combined

State with causal and non-causal parts [Khina, Kesal, Erez '11]

- $S = ({}^cS, {}^{nc}S)$
- cS – known *causally* at Tx
- ${}^{nc}S$ – known *non-causally* at Tx
- **Capacity:** $C = \max_{p(u|{}^{nc}S), x(u, {}^{nc}S, {}^cS)} [I(U; Y) - I(U; {}^{nc}S)]$
- U is independent of cS given ${}^{nc}S$

Converse

$$n(R - \epsilon_n) \stackrel{\text{Fano}}{\leq} I(W; Y_1^n) = \sum_{i=1}^n I(W; Y_i | Y_1^{i-1}) \leq \sum_{i=1}^n I(W, Y_1^{i-1}; Y_i)$$

Shannon (causal)

- $U_i \triangleq (W, Y_1^{i-1})$.
- Causality
 \Downarrow
 U_i is independent of S_i .

$$\begin{aligned} & \sum I(W, Y_1^{i-1}; Y_i) \\ &= \sum I(U_i; Y_i) \\ &\leq n \max_{p(u), p(x|u,s)} I(U; Y) \end{aligned}$$

Gel'fand-Pinsker (non-causal)

- $U_i \triangleq (W, Y_1^{i-1}, S_{i+1}^n)$
- S_i independent of (W, S_{i+1}^n)
- Chain-rule for mutual informations.

$$\begin{aligned} & \sum I(W, Y_1^{i-1}; Y_i) = \dots = \\ &= \sum I(U_i; Y_i) - I(U_i; S_i) \\ &\leq \max_{p(u|s), p(x|u,s)} \{I(U; Y) - I(U; S)\} \end{aligned}$$

Converse

In both converses...

- Achievable rate bounded by $I(U_i; Y_i) - I(U_i; S_i)$
- Auxiliary variable: $U_i = (W, Y_1^{i-1}, S_{i+1}^n)$
 - Causal case: reduces to $U_i = (W, Y_1^{i-1})$, since S_{i+1}^n is independent of Y_i
- **Can be used for the composite causal–non-causal case!**

Converse for the composite causal–non-causal case

- $S = ({}^cS, {}^{nc}S)$
- Achievable rate is bounded by $I(U_i; Y_i) - I(U_i; {}^{nc}S_i)$
- Auxiliary variable: $U_i = (W, Y_1^{i-1}, {}^{nc}S_{i+1}^n)$
- Maximization over $p(u|{}^{nc}S)$ and $x = \text{func}(u, {}^cS, {}^{nc}S)$
(U is independent of cS given ${}^{nc}S$)

Achievable

Causal (Shannon) case

- $x = t(s)$
- Strategy t is generated from W (input to equivalent channel)
- Can be thought of as “degenerated random binning” (U independent of S)

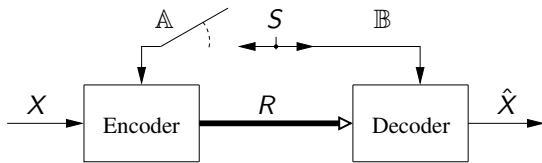
Non-causal (Gel’fand–Pinsker) case

- Use random binning w.r.t. $p(u|s)$ and $x(u, s)$

Composite causal–non-causal case

- Use random binning w.r.t. $p(u, {}^n c_S)$ and $x(u, c_S, {}^n c_S)$
- Alternatively, combine random strategies w.r.t. ${}^c S$ and random binning w.r.t. ${}^n c_S$

Source Coding with Side-information at the Decoder



- X – Source
- \hat{X} – Reconstructed (distorted) source at the decoder
- S – Side-information
- $d(X, \hat{X})$ – Distortion measure

Side-information at the encoder (\hat{A} is closed)

- S can be regarded as part of the source: $\tilde{X} \triangleq (X, S)$
- Distortion measure is w.r.t. X (and not \tilde{X})

Source Coding with Side-Information at the Decoder

Non-causal side-information [Wyner–Ziv '76]

$$R(D) = \min [I(U; X) - I(U; S)] ,$$

where minimum is over all $f : \mathcal{U} \times \mathcal{S} \rightarrow \hat{\mathcal{X}}$,
 s.t. $E [d(X, f(U, S))] \leq D$.

Causal side-information [Weissman–El Gamal '06]

$$R(D) = \min I(U; X)$$

minimum over the **same set as in the non-causal problem**.

Composite causal–non-causal side-information

$$R(D) = \min [I(U; X) - I(U; {}^{\text{nc}}S)]$$

minimum over the same set.

Separation Principle

Separation between channel and source with side-informations (SIs)
[Merhav–Shamai '03]

Separation holds between channel and source coding with

- Non-causal channel SI (Ge'fand–Pinsker)
+
Non-causal source SI (Wyner–Ziv)
- Causal channel SI (Shannon)
+
Non-causal source SI (Wyner–Ziv)

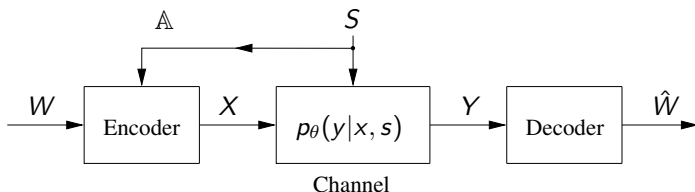
Separation between channel and source with **composite** SIs

Separation holds between channel and source coding with

Composite channel SI
+
Composite source SI

When do Both Converses Diverge?

Compound state-dependent channel:



- $\theta \in \Theta$ – “Compound parameter”: constant, unknown to Tx

Compound channel with no state knowledge (A is open)

- Worst-case capacity – Maximal rate for all θ **simultaneously**
- **Capacity:** [Blackwell et al. '59; Dobrushin '59; Wolfowitz '60]

$$C^{\text{wc}} = \max_{p(x)} \min_{\theta \in \Theta} I(X; Y)$$

When do Both Converses Diverge?

Compound Channel with state S known *causally* at Tx

Trying to generalize Shannon's converse to compound case:

- $U_i \triangleq (W, Y_1^{i-1})$
- U_i depends on the statistics of Y which is unknown!
(since θ not known)

$$n(R - \epsilon_n) \leq \sum_{i=1}^n I_{\theta}(W, Y_1^{i-1}; Y_i) \leq \sum_{i=1}^n I_{\theta}(W, \cancel{Y_1^{i-1}}, \overset{\text{causality}}{S_1^{i-1}}; Y_i)$$

- $U_i \triangleq (W, S_1^{i-1})$ – No knowledge of θ is assumed!
(Original U_i used by Shannon)
- **(Worst-case) Capacity:** [Khina, Erez '10]

$$C^{\text{WC}} = \max_{p(u), x(u, s)} \min_{\theta \in \Theta} I_{\theta}(U; Y)$$

When do Both Converses Diverge?

Compound Channel with state S known *non-causally* at Tx

- Similarly one would expect:

$$C^{\text{wc}} = \max_{p(u|s), x(u,s)} \min_{\theta \in \Theta} I_{\theta}(U; Y)$$

- This is **achievable** but **not optimal!**
 [Piantanida, Shamai '10; Nair, El Gamal, Chia '10]
 - Marton's broadcast technique exceeds this rate

Summary

- Different side-information scenarios can be treated similarly
- Recognizing these similarities allows to:
 - Solve several different scenarios at once
 - Combine results for composite/mixed scenarios
- However, in certain more complex scenarios similar/combined treatments might diverge