

The Confidential MIMO Broadcast Capacity

A Simple Derivation

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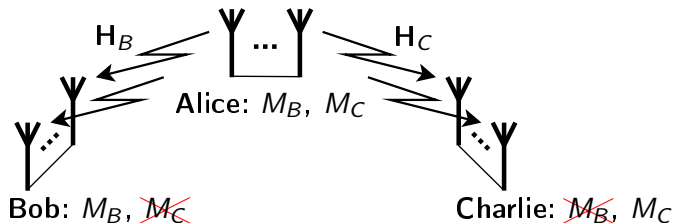
Ashish Khisti, University of Toronto

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Model: Confidential Gaussian MIMO Broadcast



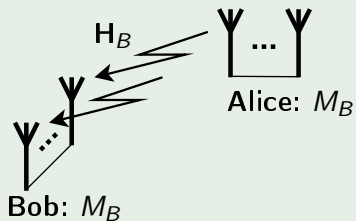
$$\mathbf{y}_B = \mathbf{H}_B \mathbf{x}_A + \mathbf{z}_B$$

$$\mathbf{y}_C = \mathbf{H}_C \mathbf{x}_A + \mathbf{z}_C$$

- \mathbf{x}_A – $N_A \times 1$ input vector
- $\mathbf{y}_B, \mathbf{y}_C$ – $N_B \times 1, N_C \times 1$ received vectors
- $\mathbf{H}_B, \mathbf{H}_C$ – $N_B \times N_A, N_C \times N_A$ channel matrices
- $\mathbf{z}_B \sim \mathcal{CN}(\mathbf{0}, \mathbf{I}_{N_B}), \mathbf{z}_C \sim \mathcal{CN}(\mathbf{0}, \mathbf{I}_{N_C})$ – noise vectors
- “Closed loop” (full channel knowledge everywhere)

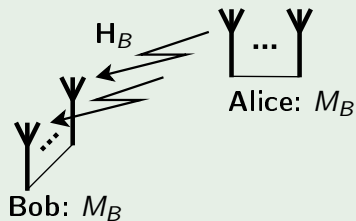
Outline of Talk (Special Cases)

MIMO Without Secrecy (No Charlie)

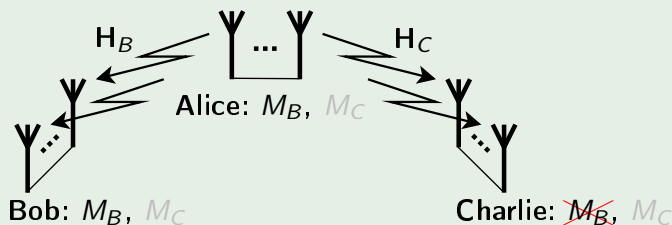


Outline of Talk (Special Cases)

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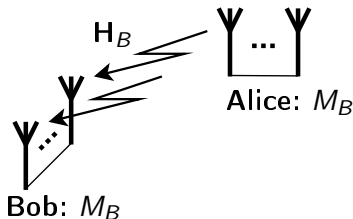


MIMO Wiretap (Maximize Rate to Bob)



MIMO Without Secrecy (No Charlie)

Practical Scheme via Matrix Decompositions



V-BLAST Scheme: QR Decomposition Based Scheme

Zero-forcing V-BLAST [Foschini '96][Wolniansky et al. '98]

- $\mathbf{H}_B = \mathbf{Q}_B \mathbf{T}_B$
- \mathbf{Q}_B – unitary; \mathbf{T}_B – triangular
- Bob applies \mathbf{Q}_B^\dagger (no SP is required by Alice)

$$\bullet \mathbf{T}_B = \begin{pmatrix} b_1 & * & * & \cdots & * \\ 0 & b_2 & * & \cdots & * \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & \cdots & b_{N_A-1} & * \\ 0 & 0 & \cdots & 0 & b_{N_A} \end{pmatrix} \Rightarrow \begin{matrix} y_1^{\text{eff}} = b_1 x_1 + z_1 \\ y_2^{\text{eff}} = b_2 x_2 + z_2 \\ \vdots \\ y_{N_A}^{\text{eff}} = b_{N_A} x_{N_A} + z_{N_A} \end{matrix}$$

- Off-diagonal elements are canceled via either:
 - Successive interference cancellation (SIC)
 - Dirty-paper coding (DPC)

V-BLAST Scheme: QR Decomposition Based Scheme

MMSE-VBLAST for a given covariance \mathbf{K} [Hassibi '00]

- $\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \mathbf{T}_B$
- \mathbf{Q}_B – unitary; $\tilde{\mathbf{Q}}_B$ – $N_B \times N_A$ submatrix of \mathbf{Q}_B
- Bob applies $\tilde{\mathbf{Q}}_B^\dagger$ (no SP is required by Alice)
- $\tilde{\mathbf{Q}}_B^\dagger$ contains Wiener-filtering (“FFE”)
- Effective noise has channel noise and “ISI” components
- Effective SNRs satisfy: $t_i^2 = 1 + \text{SNR}_i$

$$\log(t_i^2) = \log(1 + \text{SNR}_i) = I(c_i; \mathbf{y}_B | c_{i+1}^{N_A})$$
- Off-diagonal elements above diagonal canceled via SIC or DPC

V-BLAST Scheme: QR Decomposition Based Scheme

- For square invertible \mathbf{H} , ZF-VBLAST achieves:

$$R = \log \left| \mathbf{H}_B \mathbf{H}_B^\dagger \right|$$

$$\left(\text{Using } \mathbf{K} \text{ at the transmitter achieves: } R = \log \left| \mathbf{H}_B \mathbf{K} \mathbf{H}_B^\dagger \right| \right)$$

- MMSE-VBLAST achieves: $R = \log \left| \mathbf{I}_{N_B} + \mathbf{H}_B \mathbf{K} \mathbf{H}_B^\dagger \right|$

Precoded V-BLAST: Alice Applies Unitary \mathbf{V}_A

MMSE-VBLAST with precoding for a given covariance \mathbf{K}

- $\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \mathbf{T}_B$
- \mathbf{V}_A can be used to design diagonal values \Leftrightarrow design SNRs

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SVD-scheme as MMSE-VBLAST (QR)

Choosing \mathbf{V}_A of the SVD of $\mathbf{H}_B \mathbf{K}^{1/2} \Rightarrow$ SVD scheme
(no SIC/DPC needed)

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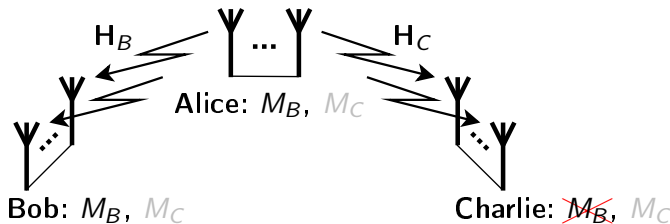
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Geometric-mean decomposition [Jiang et al. '05]/ QRS [Zhang et al. '05]

- \mathbf{V}_A is choosing s.t. all diagonal values (all SNRs) are equal
- The same codebook can be used over all subchannels
- No need for bit-loading

MIMO Wiretap (Maximize Rate to Bob)



Gaussian MIMO Wiretap: Maximize rate to Bob

Total power constraint [Khisti–Wornell '10][Oggier–Hassibi '11]

$$C_B(\mathbf{H}_B, \mathbf{H}_C) = \max_{\mathbf{K}_A: \text{tr}\{\mathbf{K}_A\} \leq P} \left[\overbrace{\log \left| \mathbf{I} + \mathbf{H}_B \mathbf{K} \mathbf{H}_B^\dagger \right|}^{I(\mathbf{X}_A; \mathbf{Y}_B)} - \overbrace{\log \left| \mathbf{I} + \mathbf{H}_C \mathbf{K} \mathbf{H}_C^\dagger \right|}^{I(\mathbf{X}_A; \mathbf{Y}_C)} \right]$$

- $X \sim \text{Gaussian}$
- Maximization over all admissible covariance matrices \mathbf{K}_A
- No explicit expression for optimal \mathbf{K}_A

Gaussian MIMO Wiretap: Maximize rate to Bob

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- $X \sim \text{Gaussian}$
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Input covariance constraint [Liu–Shamai '09]

- Replace $\text{tr}\{\mathbf{K}_A\} \leq P$ with $\mathbf{K}_A \preceq \mathbf{K}$
- Explicit exp. for optimal \mathbf{K}_A ! [Bustin–Liu–Poor–Shamai '09]

Scheme for General SNR [Kh., Kochman, Khisti ISIT'14]

$$\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \overbrace{\begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix}}^{\mathbf{T}_B}, \quad b_i^2 = 1 + \text{SNR}_i^B$$

$$\begin{bmatrix} \mathbf{H}_C \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \overbrace{\begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix}}^{\mathbf{T}_C}, \quad c_i^2 = 1 + \text{SNR}_i^E$$

- Use good SISO wiretap codes for SNR-pairs $(b_i^2 - 1, c_i^2 - 1)$
- \mathbf{V}_A of Charlie's SVD \Rightarrow Easy secrecy analysis + strong secrecy
- \mathbf{V}_A of Bob's SVD \Rightarrow No need for V-BLAST
- $\text{diag}\{\mathbf{T}_B\}, \text{diag}\{\mathbf{T}_C\}$ are const. \Rightarrow Same code over all channels

Scheme for General SNR [Kh., Kochman, Khisti ISIT'14]

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- Use good SISO wiretap codes for SNR-pairs $(b_i^2 - 1, c_i^2 - 1)$

Genie-aided secrecy-proof

- Charlie tries to recover messages sequentially (from last to first)
- For the recovery of message i all previous messages are revealed

Dirty-Paper Coding Variant of Scheme [New]

Alice:

- Cancels out off-diagonal elements of Bob via DPC
- Uses good SISO DPC wiretap codes for SNR-pairs $(b_i^2 - 1, c_i^2 - 1)$

Bob: Recovers using DPC decoder (as in MIMO without secrecy)

Genie-aided secrecy-proof

- Similar to secrecy analysis for SIC scheme
- Charlie tries to recover messages sequentially
- Messages this time are Costa-like auxiliaries (U_i)
- For the recovery of message i previous auxiliaries are revealed (from which the signals X_i can be constructed)

Reinterpretation of Optimal K_A Expression of Bustin *et al.* via GSVD

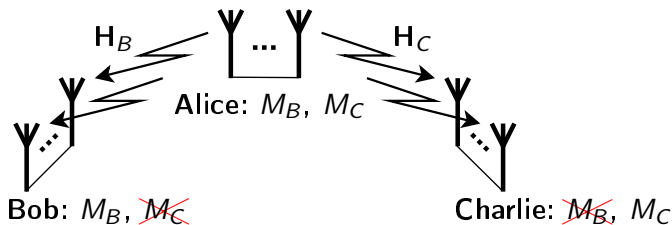
- Apply the *Generalized Singular Value Decomposition* (GSVD) to the “MMSE-VBLAST” matrices:

$$\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix} \mathbf{V}_A^\dagger$$

$$\begin{bmatrix} \mathbf{H}_C \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix} \mathbf{V}_A^\dagger$$

- $\mathbf{Q}_B, \mathbf{Q}_C, \mathbf{V}_A$ – Unitary
- $\frac{b_i}{c_i}$ decreases with i and “least balanced” [Kh.-Kochman-Erez '12]
- Choose directions corresponding to $\frac{b_i}{c_i} > 1$ and nullify the rest

Finally... Confidential MIMO Broadcast



Confidential MIMO BC Capacity [Liu–Liu–Poor–Shamai '10]

Total power constraint

Using Lemma 1 of [Weingarten–Steinberg–Shamai '06]:

$$\mathcal{C}(\mathbf{H}_B, \mathbf{H}_C, P) = \bigcup_{\mathbf{K}_A: \text{tr}\{\mathbf{K}_A\} \leq P} \left(C_B(\mathbf{H}_B, \mathbf{H}_C, \mathbf{K}_A), C_C(\mathbf{H}_B, \mathbf{H}_C, \mathbf{K}_A) \right)$$

⇒ Suffices to determine capacity region under covariance constraint

Covariance constraint

- No tension between users
- Both users achieve optimal wiretap capacities **simultaneously!**
- Rectangular capacity region

Confidential MIMO BC Capacity [Liu–Liu–Poor–Shamai '10]

Method of proof for covariance constraint

Converse: Trivial (both users achieve optimality simultaneously)

Achievable:

- Development of new achievable for MIMO wiretap channel that uses “artificial noise”
- Showing that new techniques achieves capacity (via *channel enhancement* technique)
- Double random-binning schemes and “secret dirty-paper coding”
- Alice \rightarrow Bob: Uses “standard” random binning (standard technique)
- Alice \rightarrow Charlie: Uses random binning with artificial noise (new technique)

Confidential MIMO BC Capacity [Liu–Liu–Poor–Shamai '10]

Connection to [Bustin–Liu–Poor–Shamai '09]

- Generalization of the techniques of Bustin *et al.*
- Capacity region can be represented via “generalized eigenvalues”

Confidential MIMO BC Capacity: New Derivation

- Simple derivation
- Uses standard MIMO wiretap capacity expression
- Builds upon [Bustin–Liu–Poor–Shamai '10]
- Uses standard (Costa) dirty-paper coding
- Explains why dirty-paper coding is required (for at least one of the users)
- Connection to the *generalized singular values* is natural
- Allows construction of a practical scheme

Capacity-Achieving Confidential MIMO Broadcast

$$\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix} \mathbf{v}_A^\dagger$$

$$\begin{bmatrix} \mathbf{H}_C \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix} \mathbf{v}_A^\dagger$$

- Choosing directions of $b_i > c_i$ is **optimal for Bob**
- **But...** Choosing directions of $b_i < c_i$ is **optimal for Charlie!**



Allocate $b_i > c_i$ to Bob
 Allocate $b_i < c_i$ to Charlie

Capacity-Achieving Confidential MIMO Broadcast

$$\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix} \mathbf{v}_A^\dagger$$

$$\begin{bmatrix} \mathbf{H}_C \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix} \mathbf{v}_A^\dagger$$

Scheme

Charlie: Uses SIC or DPC over elements satisfying $c_i > b_i$

Bob: Uses DPC scheme over elements satisfying $b_i > c_i$

- Alice pre-cancels interference also from Charlie's signals
- Bob can gain no info. of Charlie's message from DPC

Capacity-Achieving Confidential MIMO Broadcast

Decoupling the Modulation

- Alice can apply additional unitary block-matrix operations:

$$\tilde{\mathbf{T}}_B = \overbrace{\begin{pmatrix} \mathbf{Q}_{B;B}^\dagger & \mathbf{0} \\ \mathbf{0} & \mathbf{Q}_{B;C}^\dagger \end{pmatrix}}^{\text{Bob}} \begin{pmatrix} \mathbf{B}_B & * \\ \mathbf{0} & \mathbf{B}_C \end{pmatrix} \overbrace{\begin{pmatrix} \mathbf{V}_B & \mathbf{0} \\ \mathbf{0} & \mathbf{V}_C \end{pmatrix}}^{\text{Alice}}$$

$$\tilde{\mathbf{T}}_C = \overbrace{\begin{pmatrix} \mathbf{Q}_{C;B}^\dagger & \mathbf{0} \\ \mathbf{0} & \mathbf{Q}_{C;C}^\dagger \end{pmatrix}}^{\text{Charlie}} \begin{pmatrix} \mathbf{C}_B & * \\ \mathbf{0} & \mathbf{C}_C \end{pmatrix} \overbrace{\begin{pmatrix} \mathbf{V}_B & \mathbf{0} \\ \mathbf{0} & \mathbf{V}_C \end{pmatrix}}^{\text{Alice}}$$

- SVD of Charlie's "block" \mathbf{C}_C allows him to avoid SIC or DPC
- Other decompositions with desired properties can be applied (SVD of \mathbf{B}_C , 2-GMD of $(\mathbf{B}_B, \mathbf{B}_C)$, 2-GMD of $(\mathbf{C}_B, \mathbf{C}_C)$, etc.)

Summary

- Simple derivation of capacity region of confidential MIMO BC
- Simple extension of the result of Bustin *et al.* for MIMO WTC
- Connection to generalized singular values is immediate
- Explicitly shows why DPC by at least one user is required