Decomposing the MIMO Wiretap Channel

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Part I

Problem Setting

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Model Capacity SISO Goal

Channel Model: Gaussian MIMO Wiretap Channel



Capacity

Gaussian SISO channel capacity [Leung-Yan-Cheong, Hellman '78]

$$C_{S}(h_{B}, h_{E}) = \left[\overbrace{\log\left(1 + |h_{B}|^{2} P\right)}^{I(X;Y_{B})} - \overbrace{\log\left(1 + |h_{E}|^{2} P\right)}^{I(X;Y_{E})}\right]_{+}$$

Gaussian MIMO channel capacity [Khisti, Wornell '10][Oggier, Hassibi '11]

$$C_{S}(\mathbf{H}_{B},\mathbf{H}_{E}) = \max_{\mathbf{K}: \operatorname{trace}\{\mathbf{K}\} \leq P} \left[\overbrace{\log \left| \mathbf{I} + \mathbf{H}_{B}\mathbf{K}\mathbf{H}_{B}^{\dagger} \right|}^{I(\mathbf{X};\mathbf{Y}_{E})} - \overbrace{\log \left| \mathbf{I} + \mathbf{H}_{E}\mathbf{K}\mathbf{H}_{E}^{\dagger} \right|}^{I(\mathbf{X};\mathbf{Y}_{E})} \right]$$

Maximization over all admissible covariance matrices K

• Power constraint can be replaced with covariance constraint [Liu, Shamai '09]

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Capacity-achieving Codes for Gaussian SISO Wiretap

- Great effort in constructing practical capacity-achieving codes:
 - LDPC-based [Klinc et al. '11][Andresson, PhD '11]
 - Lattice-based [Oggier et al., submitted '13]
 - Polar-based [Mahdavifar, Vardy '11]
 - "Black-box approach" [Tyagi, Vardy, ISIT'14]
 - More...

What to do for MIMO?

Goal: Practical MIMO Scheme

Black box approach

- Signal processing (SVD-based scheme [Telatar '99], V-BLAST [Foschini '96], ...)
- Any good SISO wiretap codes
- Achieves capacity
- Gap-to-capacity dictated by gap-to-capacity of the SISO codes

Goal: Practical MIMO Scheme

Black box approach

- Signal processing (SVD-based scheme [Telatar '99], V-BLAST [Foschini '96], ...)
- Any good SISO wiretap codes
- Achieves capacity
- Gap-to-capacity dictated by gap-to-capacity of the SISO codes

High SNR [Khisti, Wornell '10]

- Generalized SVD (GSVD) + linear zero-forcing
- Any good SISO wiretap codes
- Optimal at $SNR^B, SNR^E \to \infty$

• Suboptimal at $\mathsf{SNR}^B, \mathsf{SNR}^E < \infty$

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Part II

Background: MIMO Without Secrecy

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SVD V-BLAST (QR) SVD with water-filling SVD for Given K

Practical Schemes for MIMO Without Secrecy (No Eve)

Singular-Value Decomposition (SVD) Scheme [Telatar '99]

•
$$\mathbf{H}_B = \mathbf{Q}_B \mathbf{D}_B \mathbf{V}_A^{\dagger}$$

• Alice applies V_A and Bob applies Q_B

•
$$\mathbf{D}_B = \begin{pmatrix} d_1 & 0 & 0 & \cdots & 0 \\ 0 & d_2 & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & d_{N-1} & 0 \\ 0 & \cdots & 0 & 0 & d_N \end{pmatrix} \Rightarrow \begin{array}{c} y_1 = d_1 x_1 + z_1 \\ y_2 = d_2 x_2 + z_2 \\ \Rightarrow & \vdots \\ y_N = d_N x_N + z_N \end{pmatrix}$$

- Results in parallel scalar sub-channels (each sub-channel has a different SNR)
- Apply water-filling on $\{x_1, \ldots, x_N\}$: $\mathbf{x} = \mathbf{V}_A \mathbf{W} \mathbf{c}$

SVD V-BLAST (QR) SVD with water-filling SVD for Given K

Practical Schemes for MIMO Without Secrecy (No Eve)

SVD-based scheme for a given input covariance K

•
$$\mathbf{H}_B \mathbf{K}^{1/2} = \mathbf{Q}_B \mathbf{D}_B \mathbf{V}_A^{\dagger}$$

• Alice applies $\mathbf{K}^{1/2}\mathbf{V}_A$ and Bob applies \mathbf{Q}_B

•
$$\mathbf{D}_B = \begin{pmatrix} d_1 & 0 & 0 & \cdots & 0 \\ 0 & d_2 & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & \cdots & 0 & d_{N-1} & 0 \\ 0 & \cdots & 0 & 0 & d_N \end{pmatrix} \Rightarrow \begin{cases} y_1 = d_1 x_1 + z_1 \\ y_2 = d_2 x_2 + z_2 \\ \vdots \\ y_N = d_N x_N + z_N \end{cases}$$

- Results in parallel scalar sub-channels (each sub-channel has a different SNR)
- Apply water filling on $\{x_1, \ldots, x_n\}$: $\mathbf{x} = \mathbf{V}_A \mathbf{W} \mathbf{c}$ $\mathbf{x} = \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c}$

SVD V-BLAST (QR) SVD with water-filling SVD for Given K

Practical Schemes for MIMO Without Secrecy (No Eve)

- SVD scheme with given **K** achieves : $R = \log \left| \mathbf{I}_{N_A} + \mathbf{H}_B \mathbf{K} \mathbf{H}_B^{\dagger} \right|$
- For optimal choice of K attains capacity
- Can be used to attain capacity for other covariance constraint scenarios (e.g., individual power constraints)

Practical Schemes for MIMO Without Secrecy (No Eve)

QR decomp. [Foschini'96][Wolniansky et al.'98] – zero-forcing VBLAST

- $\mathbf{H}_B = \mathbf{Q}_B \mathbf{T}_B$
- **Q**_B unitary; **T**_B triangular
- Bob applies \mathbf{Q}_B^{\dagger} (no SP is required by Alice)

•
$$\mathbf{T}_B = \begin{pmatrix} t_1 & * & * & \cdots & * \\ 0 & t_2 & * & \cdots & * \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & \cdots & t_{N-1} & * \\ 0 & 0 & \cdots & 0 & t_N \end{pmatrix} \qquad \begin{array}{c} y_1^{\text{eff}} = t_1 x_1 + z_1 \\ y_2^{\text{eff}} = t_2 x_2 + z_2 \\ \Rightarrow & \vdots \\ y_N^{\text{eff}} = t_N x_N + z_N \end{array}$$

 Off-diagonal elements are canceled via successive interference cancellation (SIC)

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SVD V-BLAST (QR) ZF MMSE Precoded

Practical Schemes for MIMO Without Secrecy (No Eve)

QRD – MMSE-VBLAST for a given covariance K [Hassibi '00]

$$\bullet \begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \mathbf{T}_B$$

- \mathbf{Q}_B unitary; $ilde{\mathbf{Q}}_B$ $N_B imes N_A$ submatrix of \mathbf{Q}_B
- Bob applies $\tilde{\mathbf{Q}}_B^{\dagger}$ (no SP is required by Alice)
- $\tilde{\mathbf{Q}}_B^{\dagger}$ contains Wiener-filtering ("FFE")
- Effective noise has channel noise and "ISI" components
- Effective SNRs satisfy: $t_i^2 = 1 + SNR_i$

$$\log(t_i^2) = \log(1 + \mathsf{SNR}_i) = I(c_i; \mathbf{y}_B | c_{i+1}^{N_A})$$

Off-diagonal elements above diagonal canceled via SIC

SVD V-BLAST (QR) ZF MMSE Precoded

Practical Schemes for MIMO Without Secrecy (No Eve)

- For square invertible **H**, ZF-VBLAST achieves: $R = |\mathbf{H}_B \mathbf{H}_B^{\dagger}|$ (Using **K** at the transmitter achieves: $R = |\mathbf{H}_B \mathbf{K} \mathbf{H}_B^{\dagger}|$)
- MMSE-VBLAST achieves: $R = \left| \mathbf{I}_{N_B} + \mathbf{H}_B \mathbf{K} \mathbf{H}_B^{\dagger} \right|$



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SVD V-BLAST (QR) ZF MMSE Precoded

Practical Schemes for MIMO Without Secrecy (No Eve)

MMSE-VBLAST with precoding for a given covariance K

$$\mathbf{P} \begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \mathbf{T}_B$$

• V_A can be used to design diagonal values \Leftrightarrow design SNRs

SVD-scheme as MMSE-VBLAST (QR)

Choosing V_A of the SVD of $H_B K^{1/2} \Rightarrow$ SVD scheme (no SIC needed) SVD V-BLAST (QR) ZF MMSE Precoded

Practical Schemes for MIMO Without Secrecy (No Eve)

MMSE-VBLAST with precoding for a given covariance K

$$\mathbf{P} \begin{bmatrix} \mathbf{H}_{B} \mathbf{K}^{1/2} \mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{B} \mathbf{T}_{B}$$

• V_A can be used to design diagonal values \Leftrightarrow design SNRs

SVD-scheme as MMSE-VBLAST (QR)

Choosing
$$\mathbf{V}_A$$
 of the SVD of $\mathbf{H}_B \mathbf{K}^{1/2} \Rightarrow$ SVD scheme
(no SIC needed)

Geometric-mean decomposition [Jiang et al. '05]/ QRS [Zhang et al. '05]

- V_A is choosing s.t. all diagonal values (all SNRs) are equal
- The same codebook can be used over all subchannels
- No need for bit-loading

Part III

Back to Wiretap...

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New Scheme for Finite SNR



• Use any good SISO wiretap codes for SNR-pairs $(b_i^2 - 1, e_i^2 - 1)$

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New Scheme for Finite SNR

$$\begin{bmatrix} \mathbf{H}_{B}\mathbf{K}^{1/2}\mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{B}\overbrace{\begin{pmatrix} b_{1} & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N} \end{pmatrix}}^{\mathbf{T}_{B}}, \quad b_{i}^{2} = 1 + \mathrm{SNR}_{i}^{B}$$
$$\begin{bmatrix} \mathbf{H}_{E}\mathbf{K}^{1/2}\mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{E}\overbrace{\begin{pmatrix} e_{1} & * & * \\ 0 & \ddots & * \\ 0 & 0 & e_{N} \end{pmatrix}}^{\mathbf{T}_{E}}, \quad e_{i}^{2} = 1 + \mathrm{SNR}_{i}^{E}$$

• Use any good SISO wiretap codes for SNR-pairs $(b_i^2 - 1, e_i^2 - 1)$

• V_A of Eve's SVD \Rightarrow Easy secrecy analysis + strong secrecy

New Scheme for Finite SNR

$$\begin{bmatrix} \mathbf{H}_{B}\mathbf{K}^{1/2}\mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{B}\overbrace{\begin{pmatrix} b_{1} & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N} \end{pmatrix}}^{\mathbf{T}_{B}}, \quad b_{i}^{2} = 1 + \mathrm{SNR}_{i}^{B}$$
$$\begin{bmatrix} \mathbf{H}_{E}\mathbf{K}^{1/2}\mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{E}\overbrace{\begin{pmatrix} e_{1} & * & * \\ 0 & \ddots & * \\ 0 & 0 & e_{N} \end{pmatrix}}^{\mathbf{T}_{E}}, \quad e_{i}^{2} = 1 + \mathrm{SNR}_{i}^{E}$$

• Use any good SISO wiretap codes for SNR-pairs $(b_i^2 - 1, e_i^2 - 1)$

• V_A of Eve's SVD \Rightarrow Easy secrecy analysis + strong secrecy

• V_A of Bob's SVD \Rightarrow No need for V-BLAST

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New Scheme for Finite SNR

$$\begin{bmatrix} \mathbf{H}_{B}\mathbf{K}^{1/2}\mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{B}\overbrace{\begin{pmatrix} b_{1} & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N} \end{pmatrix}}^{\mathbf{T}_{B}}, \quad b_{i}^{2} = 1 + \mathrm{SNR}_{i}^{B}$$
$$\begin{bmatrix} \mathbf{H}_{E}\mathbf{K}^{1/2}\mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{E}\overbrace{\begin{pmatrix} e_{1} & * & * \\ 0 & \ddots & * \\ 0 & 0 & e_{N} \end{pmatrix}}^{\mathbf{T}_{E}}, \quad e_{i}^{2} = 1 + \mathrm{SNR}_{i}^{E}$$

• Use any good SISO wiretap codes for SNR-pairs $(b_i^2 - 1, e_i^2 - 1)$

• V_A of Eve's SVD \Rightarrow Easy secrecy analysis + strong secrecy

- V_A of Bob's SVD \Rightarrow No need for V-BLAST
- diag{ T_B }, diag{ T_E } are const. \Rightarrow Same code over all channels

 ${\ensuremath{\, \circ }}$ Take optimal covariance matrix ${\ensuremath{K}}$

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- Take optimal covariance matrix K
- Use optimal SVD-based scheme matched for H_E with K:

$$\mathbf{H}_{E}\mathbf{K}^{1/2} = \mathbf{Q}_{E}\mathbf{D}_{E}\mathbf{V}_{A}^{\dagger}$$

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- Take optimal covariance matrix K
- Use optimal SVD-based scheme matched for H_E with K:

$$\mathbf{H}_{E}\mathbf{K}^{1/2}=\mathbf{Q}_{E}\mathbf{D}_{E}\mathbf{V}_{A}^{\dagger}$$

• Alice: $\mathbf{x} = \mathbf{K}^{1/2} \mathbf{V}_{A} \mathbf{c}$

- Take optimal covariance matrix K
- Use optimal SVD-based scheme matched for H_E with K:

$$\mathbf{H}_{E}\mathbf{K}^{1/2} = \mathbf{Q}_{E}\mathbf{D}_{E}\mathbf{V}_{A}^{\dagger}$$

- Alice: $\mathbf{x} = \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c}$
- Eve: $\tilde{\mathbf{y}}_E = \mathbf{Q}_E^{\dagger} \mathbf{y}_E = \mathbf{Q}_E^{\dagger} \mathbf{H}_E \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c} \mathbf{Q}_E^{\dagger} + \mathbf{z}_E = \mathbf{D}_E \mathbf{c} + \tilde{\mathbf{z}}_E$

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- Take optimal covariance matrix K
- Use optimal SVD-based scheme matched for H_E with K:

$$\mathbf{H}_{E}\mathbf{K}^{1/2}=\mathbf{Q}_{E}\mathbf{D}_{E}\mathbf{V}_{A}^{\dagger}$$

- Alice: $\mathbf{x} = \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c}$
- Eve: $\tilde{\mathbf{y}}_E = \mathbf{Q}_E^{\dagger} \mathbf{y}_E = \mathbf{Q}_E^{\dagger} \mathbf{H}_E \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c} \mathbf{Q}_E^{\dagger} + \mathbf{z}_E = \mathbf{D}_E \mathbf{c} + \tilde{\mathbf{z}}_E$
- Eve sees parallel subchannels \Rightarrow Easy secrecy analysis

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- Take optimal covariance matrix K
- Use optimal SVD-based scheme matched for \mathbf{H}_E with \mathbf{K} :

$$\mathbf{H}_{E}\mathbf{K}^{1/2}=\mathbf{Q}_{E}\mathbf{D}_{E}\mathbf{V}_{A}^{\dagger}$$

- Alice: $\mathbf{x} = \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c}$
- Eve: $\tilde{\mathbf{y}}_E = \mathbf{Q}_E^{\dagger} \mathbf{y}_E = \mathbf{Q}_E^{\dagger} \mathbf{H}_E \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c} \mathbf{Q}_E^{\dagger} + \mathbf{z}_E = \mathbf{D}_E \mathbf{c} + \tilde{\mathbf{z}}_E$
- Eve sees parallel subchannels \Rightarrow Easy secrecy analysis
- Bob Uses MMSE-VBLAST: $\tilde{y}_B = \tilde{Q}_B y_B = \tilde{T}_B c + \tilde{z}_B$

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- Take optimal covariance matrix K
- Use optimal SVD-based scheme matched for H_E with K:

$$\mathbf{H}_{E}\mathbf{K}^{1/2}=\mathbf{Q}_{E}\mathbf{D}_{E}\mathbf{V}_{A}^{\dagger}$$

- Alice: $\mathbf{x} = \mathbf{K}^{1/2} \mathbf{V}_{A} \mathbf{c}$
- Eve: $\tilde{\mathbf{y}}_E = \mathbf{Q}_E^{\dagger} \mathbf{y}_E = \mathbf{Q}_E^{\dagger} \mathbf{H}_E \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c} \mathbf{Q}_E^{\dagger} + \mathbf{z}_E = \mathbf{D}_E \mathbf{c} + \tilde{\mathbf{z}}_E$
- Eve sees parallel subchannels \Rightarrow Easy secrecy analysis
- Bob Uses MMSE-VBLAST: $\tilde{y}_B = \tilde{Q}_B y_B = \tilde{T}_B c + \tilde{z}_B$
- Use SISO wiretap codes over parallel wiretap subchannels: $(SNR_i^B, SNR_i^E) \leftrightarrow (\tilde{t}_i^B, d_i^E)$

- Take optimal covariance matrix K
- Use optimal SVD-based scheme matched for \mathbf{H}_E with \mathbf{K} :

$$\mathbf{H}_{E}\mathbf{K}^{1/2}=\mathbf{Q}_{E}\mathbf{D}_{E}\mathbf{V}_{A}^{\dagger}$$

- Alice: $\mathbf{x} = \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c}$
- Eve: $\tilde{\mathbf{y}}_E = \mathbf{Q}_E^{\dagger} \mathbf{y}_E = \mathbf{Q}_E^{\dagger} \mathbf{H}_E \mathbf{K}^{1/2} \mathbf{V}_A \mathbf{c} \mathbf{Q}_E^{\dagger} + \mathbf{z}_E = \mathbf{D}_E \mathbf{c} + \tilde{\mathbf{z}}_E$
- Eve sees parallel subchannels \Rightarrow Easy secrecy analysis
- Bob Uses MMSE-VBLAST: $\tilde{y}_B = \tilde{Q}_B y_B = \tilde{T}_B c + \tilde{z}_B$
- Use SISO wiretap codes over parallel wiretap subchannels: $(SNR_i^B, SNR_i^E) \leftrightarrow (\tilde{t}_i^B, d_i^E)$
- Achieves capacity for optimal K:

$$R = \underbrace{\left| \mathbf{I}_{N_A} + \mathbf{H}_B \mathbf{K} \mathbf{H}_B^{\dagger} \right|}_{\text{MMSE-VBLAST @ Bob}} - \underbrace{\left| \mathbf{I}_{N_A} + \mathbf{H}_E \mathbf{K} \mathbf{H}_E^{\dagger} \right|}_{\text{SVD @ Eve}} \equiv C_S$$

- SVD of Eve's channel allows easy secrecy-constraints proof
- Strong/weak secrecy of the SISO codes ↓ Strong/weak secrecy of the MIMO scheme

Can we use SVD for Bob's channel?

Can the same SISO codebook be used over all subchannels?

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Superposition Coding for the Memoryless Wiretap Channel

Theorem

- $p(y_B|x)$ and $p(y_E|x)$ transition distributions for Bob and Eve
- $x = \varphi(c_1, \ldots, c_{N_A})$
- c_i drawn i.i.d. $p_{c_i}(\cdot)$
- Weak-secrecy rates of

$$R_{i} = I\left(c_{i}; y_{B} \middle| c_{i+1}^{N_{A}}\right) - I\left(c_{i}; y_{E} \middle| c_{i+1}^{N_{A}}\right), \quad i = 1, \dots, N_{A}$$

are achievable

Genie-aided secrecy-proof

- Eve tries to recover sub-messages sequentially (from last to first)
- For the recovery of sub-message i all previous sub-messages $(i = i + 1, ..., N_A)$ are revealed

General Multi-Stream Scheme for the Gaussian MIMO Wiretap

• Apply QR decompositions for both Bob and Eve:

$$\begin{bmatrix} \mathbf{H}_{B}\mathbf{K}^{1/2}\mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{B}\mathbf{T}_{B}, \qquad \begin{bmatrix} \mathbf{H}_{E}\mathbf{K}^{1/2}\mathbf{V}_{A} \\ \mathbf{I}_{N_{A}} \end{bmatrix} = \mathbf{Q}_{E}\mathbf{T}_{E}$$

• The resulting SNRs satisfy:

$$log(1 + SNR_{B;i}) = log(b_i^2) = I(c_i; \mathbf{y}_B | c_{i+1}^{N_A}) log(1 + SNR_{E;i}) = log(e_i^2) = I(c_i; \mathbf{y}_E | c_{i+1}^{N_A})$$

• Superposition theorem \Rightarrow Capacity is achieved with SISO codebooks of SNR-pairs $(b_i^2 - 1, e_i^2 - 1)$

\implies Alice can apply any V_A!

General Multi-Stream Scheme for the Gaussian MIMO Wiretap

SVD @ Bob

Choose V_A such that

$$\mathbf{H}_B\mathbf{K}^{1/2}=\mathbf{Q}_B\mathbf{D}_B\mathbf{V}_A^\dagger$$

• Bob: $\tilde{\mathbf{y}}_B = \mathbf{D}_B \mathbf{c} + \tilde{\mathbf{z}}_B$

Same codebook over all subchannels

• Choose V_A s.t. T_B and T_E have constant diagonals:

$$SNR_1^B = SNR_2^B = \cdots = SNR_N^B$$

 $SNR_1^E = SNR_2^E = \cdots = SNR_N^E$

• Possible using a space-time structure [Khina et al. '11]

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Part IV

Supplementary

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• Apply Generalized SVD (GSVD) to \mathbf{H}_B and \mathbf{H}_E :

 $\mathbf{H}_B = \mathbf{Q}_B \mathbf{D}_B \mathbf{A}$ $\mathbf{H}_E = \mathbf{Q}_E \mathbf{D}_E \mathbf{A}$

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• Apply Generalized SVD (GSVD) to \mathbf{H}_B and \mathbf{H}_E :

• \mathbf{A} – invertible matrix; \mathbf{Q}_B , \mathbf{Q}_E – unitary

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• Apply Generalized SVD (GSVD) to H_B and H_E :

 $\mathbf{H}_B = \mathbf{Q}_B \mathbf{D}_B \mathbf{A}$ $\mathbf{H}_E = \mathbf{Q}_E \mathbf{D}_E \mathbf{A}$

- A invertible matrix; Q_B, Q_E unitary
- \mathbf{D}_B , \mathbf{D}_E diagonal, s.t. $\mathbf{D}_B^2 + \mathbf{D}_E^2 = \mathbf{I}$

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• Apply Generalized SVD (GSVD) to \mathbf{H}_B and \mathbf{H}_E :

 $\mathbf{H}_B = \mathbf{Q}_B \mathbf{D}_B \mathbf{A}$ $\mathbf{H}_E = \mathbf{Q}_E \mathbf{D}_E \mathbf{A}$

- A invertible matrix; Q_B, Q_E unitary
- \mathbf{D}_B , \mathbf{D}_E diagonal, s.t. $\mathbf{D}_B^2 + \mathbf{D}_E^2 = \mathbf{I}$
- Alice uses linear zero-forcing inverts **A** @ Tx: $\mathbf{x} = \mathbf{A}^{-1}\mathbf{c}$

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• Apply Generalized SVD (GSVD) to \mathbf{H}_B and \mathbf{H}_E :

 $\mathbf{H}_B = \mathbf{Q}_B \mathbf{D}_B \mathbf{A}$ $\mathbf{H}_E = \mathbf{Q}_E \mathbf{D}_E \mathbf{A}$

- A invertible matrix; Q_B, Q_E unitary
- \mathbf{D}_B , \mathbf{D}_E diagonal, s.t. $\mathbf{D}_B^2 + \mathbf{D}_E^2 = \mathbf{I}$
- Alice uses linear zero-forcing inverts **A** @ Tx: $\mathbf{x} = \mathbf{A}^{-1}\mathbf{c}$
- \bullet Bob and Eve apply ${\bf Q}_B^\dagger$ and ${\bf Q}_E^\dagger$

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• Apply Generalized SVD (GSVD) to H_B and H_E :

- \mathbf{A} invertible matrix; \mathbf{Q}_B , \mathbf{Q}_E unitary
- \mathbf{D}_B , \mathbf{D}_E diagonal, s.t. $\mathbf{D}_B^2 + \mathbf{D}_E^2 = \mathbf{I}$
- Alice uses linear zero-forcing inverts **A** @ Tx: $\mathbf{x} = \mathbf{A}^{-1}\mathbf{c}$
- \bullet Bob and Eve apply ${\bf Q}_B^\dagger$ and ${\bf Q}_E^\dagger$
- Results in parallel subchannels: **D**_B, **D**_E

• Apply Generalized SVD (GSVD) to H_B and H_E :

 $\mathbf{H}_B = \mathbf{Q}_B \mathbf{D}_B \mathbf{A}$ $\mathbf{H}_E = \mathbf{Q}_E \mathbf{D}_E \mathbf{A}$

- \mathbf{A} invertible matrix; \mathbf{Q}_B , \mathbf{Q}_E unitary
- \mathbf{D}_B , \mathbf{D}_E diagonal, s.t. $\mathbf{D}_B^2 + \mathbf{D}_E^2 = \mathbf{I}$
- Alice uses linear zero-forcing inverts **A** @ Tx: $\mathbf{x} = \mathbf{A}^{-1}\mathbf{c}$
- \bullet Bob and Eve apply ${\bf Q}_B^\dagger$ and ${\bf Q}_E^\dagger$
- Results in parallel subchannels: **D**_B, **D**_E
- Transmit only over subchannels with d_{B;i} > d_{E;i} using good SISO wiretap codes

@ High-SNR $(P \rightarrow \infty)$

- Approaches capacity: $\sum_{i=1}^{N_A} \log \frac{1 + Pd_{B_{ii}}^2}{1 + Pd_{E_{ii}}^2} \approx \sum_{i=1}^{N_A} \log \frac{Pd_{B_{ii}}^2}{Pd_{E_{ii}}^2}$
- Linear zero-forcing attains capacity at high SNR (In contrast to communication without secrecy!)
- Strong/weak secrecy of the SISO codes ↓ Strong/weak secrecy of the MIMO scheme
- Wasteful at finite SNR! ^(C)

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