

# The Confidential MIMO Broadcast Capacity

## A Simple Derivation

Anatoly Khina, Tel Aviv University

Joint work with:

Yuval Kochman, Hebrew University

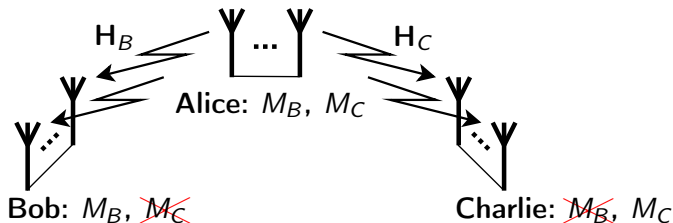
Ashish Khisti, University of Toronto

ISIT 2015

Hong Kong

June 18, 2015

## Model: Confidential Gaussian MIMO Broadcast



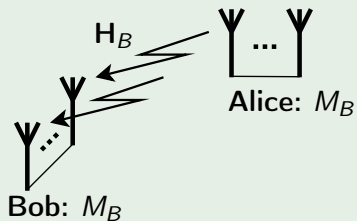
$$\mathbf{y}_B = \mathbf{H}_B \mathbf{x}_A + \mathbf{z}_B$$

$$\mathbf{y}_C = \mathbf{H}_C \mathbf{x}_A + \mathbf{z}_C$$

- $\mathbf{x}_A$  –  $N_A \times 1$  input vector
- $\mathbf{y}_B, \mathbf{y}_C$  –  $N_B \times 1, N_C \times 1$  received vectors
- $\mathbf{H}_B, \mathbf{H}_C$  –  $N_B \times N_A, N_C \times N_A$  channel matrices
- $\mathbf{z}_B \sim \mathcal{CN}(\mathbf{0}, \mathbf{I}_{N_B}), \mathbf{z}_C \sim \mathcal{CN}(\mathbf{0}, \mathbf{I}_{N_C})$  – noise vectors
- “Closed loop” (full channel knowledge everywhere)

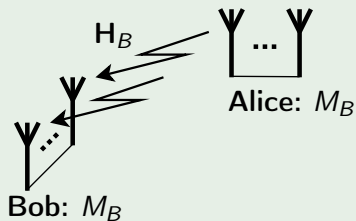
## Outline of Talk (Special Cases)

## MIMO Without Secrecy (No Charlie)

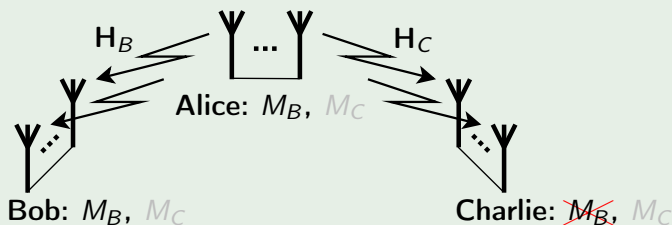


# Outline of Talk (Special Cases)

## MIMO Without Secrecy (No Charlie)

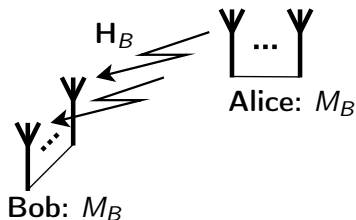


## MIMO Wiretap (Maximize Rate to Bob)



# MIMO Without Secrecy (No Charlie)

## Practical Scheme via Matrix Decompositions



# MIMO Without Secrecy (No Charlie)

## Total power constraint: Capacity

$$C(\mathbf{H}_B) = \max_{\mathbf{K}: \text{tr}\{\mathbf{K}\} \leq P} \overbrace{\log |I + \mathbf{H}_B \mathbf{K} \mathbf{H}_B|}^{I(\mathbf{X}_A; \mathbf{Y}_B)}$$

- $X \sim \text{Gaussian}$
- Maximization over all admissible covariance matrices  $\mathbf{K}$
- Explicit expression for optimal  $\mathbf{K}$  via SVD [Telatar '99]

## Input covariance constraint: Capacity

- Replace  $\text{tr}\{\mathbf{K}\} \leq P$  with  $\mathbf{K} \preceq \bar{\mathbf{K}}$
- Use all of the “available covariance”  $\bar{\mathbf{K}}$

# V-BLAST Scheme: QR Decomposition Based Scheme

## Zero-forcing V-BLAST [Foschini '96][Wolniansky et al. '98]

- $\mathbf{H}_B = \mathbf{Q}_B \mathbf{T}_B$
- $\mathbf{Q}_B$  – unitary;  $\mathbf{T}_B$  – triangular
- Bob applies  $\mathbf{Q}_B^\dagger$  (no SP is required by Alice)

$$\bullet \mathbf{T}_B = \begin{pmatrix} b_1 & * & * & \cdots & * \\ 0 & b_2 & * & \cdots & * \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & \cdots & b_{N_A-1} & * \\ 0 & 0 & \cdots & 0 & b_{N_A} \end{pmatrix} \Rightarrow \begin{aligned} y_1^{\text{eff}} &= b_1 \tilde{x}_1 + z_1 \\ y_2^{\text{eff}} &= b_2 \tilde{x}_2 + z_2 \\ &\vdots \\ y_{N_A}^{\text{eff}} &= b_{N_A} \tilde{x}_{N_A} + z_{N_A} \end{aligned}$$

- Off-diagonal elements are canceled via either:
  - Successive interference cancellation (SIC)
  - Dirty-paper coding (DPC)

# V-BLAST Scheme: QR Decomposition Based Scheme

## MMSE-VBLAST for a given covariance $\mathbf{K}$ [Hassibi '00]

- $\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \mathbf{T}_B$
- $\mathbf{Q}_B$  – unitary;  $\tilde{\mathbf{Q}}_B$  –  $N_B \times N_A$  submatrix of  $\mathbf{Q}_B$
- Bob applies  $\tilde{\mathbf{Q}}_B^\dagger$  (no SP is required by Alice)
- $\tilde{\mathbf{Q}}_B^\dagger$  contains Wiener-filtering (“FFE”)
- Effective noise has channel noise and “ISI” components
- Effective SNRs satisfy:  $b_i^2 = 1 + \text{SNR}_i$   

$$\log(b_i^2) = \log(1 + \text{SNR}_i) = I(\tilde{x}_i; \mathbf{y}_B | \tilde{x}_{i+1}^{N_A})$$
- Off-diagonal elements above diagonal canceled via SIC or DPC



# V-BLAST Scheme: QR Decomposition Based Scheme

- For square invertible  $\mathbf{H}$ , ZF-VBLAST achieves:

$$R = \log \left| \mathbf{H}_B \mathbf{H}_B^\dagger \right|$$

$$\left( \text{Using } \mathbf{K} \text{ at the transmitter achieves: } R = \log \left| \mathbf{H}_B \mathbf{K} \mathbf{H}_B^\dagger \right| \right)$$

- MMSE-VBLAST achieves:

$$R = \sum_{i=1}^{N_A} \log b_i^2 = \log \left| \mathbf{I}_{N_B} + \mathbf{H}_B \mathbf{K} \mathbf{H}_B^\dagger \right|$$

Precoded V-BLAST: Alice Applies Unitary  $\mathbf{V}_A$ MMSE-VBLAST with precoding for a given covariance  $\mathbf{K}$ 

- $\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \mathbf{T}_B$
- $\mathbf{V}_A$  can be used to design diagonal values  $\Leftrightarrow$  design SNRs

Precoded V-BLAST: Alice Applies Unitary  $\mathbf{V}_A$ MMSE-VBLAST with precoding for a given covariance  $\mathbf{K}$ 

- $\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \mathbf{T}_B$
- $\mathbf{V}_A$  can be used to design diagonal values  $\Leftrightarrow$  design SNRs

## SVD-scheme as MMSE-VBLAST (QR)

Choosing  $\mathbf{V}_A$  of the SVD of  $\mathbf{H}_B \mathbf{K}^{1/2} \Rightarrow$  SVD scheme  
(no SIC/DPC needed)

# Precoded V-BLAST: Alice Applies Unitary $\mathbf{V}_A$

## MMSE-VBLAST with precoding for a given covariance $\mathbf{K}$

- $$\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \mathbf{T}_B$$
- $\mathbf{V}_A$  can be used to design diagonal values  $\Leftrightarrow$  design SNRs

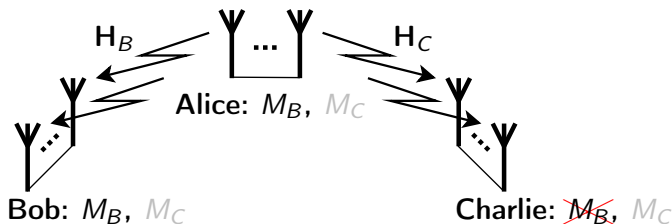
## SVD-scheme as MMSE-VBLAST (QR)

Choosing  $\mathbf{V}_A$  of the SVD of  $\mathbf{H}_B \mathbf{K}^{1/2} \Rightarrow$  SVD scheme  
(no SIC/DPC needed)

## Geometric-mean decomposition [Jiang et al. '05]/ QRS [Zhang et al. '05]

- $\mathbf{V}_A$  is choosing s.t. all diagonal values (all SNRs) are equal
- The same codebook can be used over all subchannels
- No need for bit-loading

## MIMO Wiretap (Maximize Rate to Bob)



## Gaussian MIMO Wiretap: Maximize rate to Bob

Total power constraint [Khisti–Wornell '10][Oggier–Hassibi '11]

$$C_S(\mathbf{H}_B, \mathbf{H}_C) = \max_{\mathbf{K}: \text{tr}\{\mathbf{K}\} \leq P} \left[ \overbrace{\log \left| \mathbf{I} + \mathbf{H}_B \mathbf{K} \mathbf{H}_B^\dagger \right|}^{I(\mathbf{X}_A; \mathbf{Y}_B)} - \overbrace{\log \left| \mathbf{I} + \mathbf{H}_C \mathbf{K} \mathbf{H}_C^\dagger \right|}^{I(\mathbf{X}_A; \mathbf{Y}_C)} \right]$$

- $X \sim \text{Gaussian}$
- Maximization over all admissible covariance matrices  $\mathbf{K}$
- No explicit expression for optimal  $\mathbf{K}$

# Gaussian MIMO Wiretap: Maximize rate to Bob

Total power constraint [Khisti–Wornell '10][Oggier–Hassibi '11]

$$C_S(\mathbf{H}_B, \mathbf{H}_C) = \max_{\mathbf{K}: \text{tr}\{\mathbf{K}\} \leq P} \left[ \overbrace{\log \left| \mathbf{I} + \mathbf{H}_B \mathbf{K} \mathbf{H}_B^\dagger \right|}^{I(\mathbf{X}_A; \mathbf{Y}_B)} - \overbrace{\log \left| \mathbf{I} + \mathbf{H}_C \mathbf{K} \mathbf{H}_C^\dagger \right|}^{I(\mathbf{X}_A; \mathbf{Y}_C)} \right]$$

- $X \sim \text{Gaussian}$
- Maximization over all admissible covariance matrices  $\mathbf{K}$
- No explicit expression for optimal  $\mathbf{K}$

Input covariance constraint [Liu–Shamai '09]

- Replace  $\text{tr}\{\mathbf{K}\} \leq P$  with  $\mathbf{K} \preceq \bar{\mathbf{K}}$
- Explicit exp. for optimal  $\mathbf{K}$ ! [Bustin–Liu–Poor–Shamai '09]

## Scheme for General SNR [Kh., Kochman, Khisti ISIT'14]

$$\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \overbrace{\begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix}}^{\mathbf{T}_B}, \quad b_i^2 = 1 + \text{SNR}_i^B$$

$$\begin{bmatrix} \mathbf{H}_C \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \overbrace{\begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix}}^{\mathbf{T}_C}, \quad c_i^2 = 1 + \text{SNR}_i^E$$

- Use good SISO wiretap codes for SNR-pairs  $(b_i^2 - 1, c_i^2 - 1)$
- $\mathbf{V}_A$  of Charlie's SVD  $\Rightarrow$  Easy secrecy analysis + strong secrecy
- $\mathbf{V}_A$  of Bob's SVD  $\Rightarrow$  No need for V-BLAST
- $\text{diag}\{\mathbf{T}_B\}, \text{diag}\{\mathbf{T}_C\}$  are const.  $\Rightarrow$  Same code over all channels



## Scheme for General SNR [Kh., Kochman, Khisti ISIT'14]

$$\begin{bmatrix} \mathbf{H}_B \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \overbrace{\begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix}}^{\mathbf{T}_B}, \quad b_i^2 = 1 + \text{SNR}_i^B$$

$$\begin{bmatrix} \mathbf{H}_C \mathbf{K}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \overbrace{\begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix}}^{\mathbf{T}_C}, \quad c_i^2 = 1 + \text{SNR}_i^C$$

- Use good SISO wiretap codes for SNR-pairs  $(b_i^2 - 1, c_i^2 - 1)$

## Genie-aided secrecy-proof

- Charlie tries to recover messages sequentially (from last to first)
- For the recovery of message  $i$  all previous messages are revealed

# Dirty-Paper Coding Variant of Scheme [New]

## Alice:

- Cancels out off-diagonal elements of Bob via DPC
- Uses good SISO DPC wiretap codes for SNR-pairs  $(b_i^2 - 1, c_i^2 - 1)$

**Bob:** Recovers using DPC decoder (as in MIMO without secrecy)

## Genie-aided secrecy-proof

- Similar to secrecy analysis for SIC scheme
- Charlie tries to recover messages sequentially
- Messages this time are Costa-like auxiliaries  $(U_i)$
- For the recovery of message  $i$  previous auxiliaries are revealed (from which the signals  $X_i$  can be constructed)

Restatement of Optimal  $K$  Expression of Bustin *et al.* via GSVDMIMO Wiretap Capacity under Covariance Constraint  $\mathbf{K} \preceq \bar{\mathbf{K}}$ 

- Apply  $\mathbf{V}_A$  of the *Generalized Singular Value Decomposition* (GSVD) to the “MMSE–VBLAST” matrices:

$$\begin{bmatrix} \mathbf{H}_B \bar{\mathbf{K}}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix}$$

$$\begin{bmatrix} \mathbf{H}_C \bar{\mathbf{K}}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix}$$

- $C_S(\mathbf{H}_B, \mathbf{H}_C) = \sum_{i=1}^{N_A} \left[ \log \frac{b_i}{c_i} \right]^+$
- $\mathbf{K}_{\text{opt}} = \bar{\mathbf{K}}^{1/2} \mathbf{V}_A \mathbf{I}_B \mathbf{V}_A^\dagger \bar{\mathbf{K}}^{1/2 \dagger}$
- $\mathbf{I}_B$  – Diagonal matrix with 1 whenever  $b_i > c_i$  and 0 otherwise

## Wiretap under Cov. Constraint: Alternative Proof Outline

- W.l.o.g.,  $\mathbf{K} \preceq \bar{\mathbf{K}}$  can be written as  $\mathbf{K} = \bar{\mathbf{K}}^{1/2} \mathbf{V}_A \mathbf{D} \mathbf{V}_A^\dagger \bar{\mathbf{K}}^{\dagger/2}$
- $\mathbf{D}$  is non-negative diagonal with  $0 \leq D_{ii} \leq 1$
- For any  $\mathbf{V}_A$ :  $I(\mathbf{H}_B, \mathbf{K}) - I(\mathbf{H}_C, \mathbf{K}) = \sum_{i=1}^N \log \frac{b_i^2}{c_i^2}$
- Optimal  $\mathbf{D}$  for a given  $\mathbf{V}_A$ : truncation

$$C_S(\mathbf{H}_B, \mathbf{H}_C, \bar{\mathbf{K}}) = \max_{\mathbf{V}_A} \sum_{i=1}^N \left[ \log \frac{b_i^2}{c_i^2} \right]^+$$

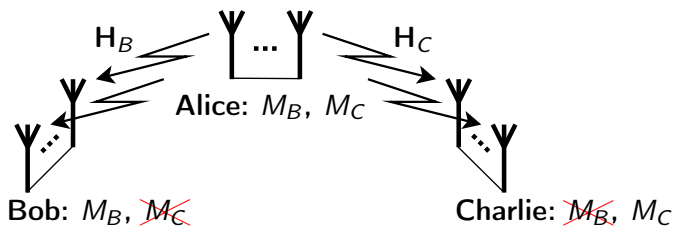
- By multiplicative majorization of joint unitary triangularization [Khina, Kochman, Erez SP'12],  $\mathbf{V}_A$  of the GSVD is optimal

## Optimal covariance matrix

$$\mathbf{K}_{\text{opt}} = \bar{\mathbf{K}}^{1/2} \mathbf{V}_A \mathbf{I}_B \mathbf{V}_A^\dagger \bar{\mathbf{K}}^{1/2 \dagger}$$

- $\mathbf{I}_B$  – Diagonal matrix with 1 whenever  $b_i > c_i$  and 0 otherwise

# Finally... Confidential MIMO Broadcast



## Confidential MIMO BC Capacity [Liu–Liu–Poor–Shamai '10]

## Total power constraint

Using Lemma 1 of [Weingarten–Steinberg–Shamai '06]:

$$\mathcal{C}(\mathbf{H}_B, \mathbf{H}_C, P) = \bigcup_{\mathbf{K}: \text{tr}\{\mathbf{K}\}=P} \left( C_S(\mathbf{H}_B, \mathbf{H}_C, \mathbf{K}), C_S(\mathbf{H}_C, \mathbf{H}_B, \mathbf{K}) \right)$$

⇒ Suffices to determine capacity region under covariance constraint

## Covariance constraint

- No tension between users
- Both users achieve optimal wiretap capacities **simultaneously!**
- Rectangular capacity region

## Confidential MIMO BC Capacity [Liu–Liu–Poor–Shamai '10]

## Method of proof for covariance constraint

**Converse:** Trivial (both users achieve optimality simultaneously)

**Achievable:**

- Development of new achievable for MIMO wiretap channel that uses “artificial noise”
- Showing that new techniques achieves capacity (via *channel enhancement* technique)
- Double random-binning schemes and “secret dirty-paper coding”
- Alice  $\rightarrow$  Bob: Uses “standard” random binning (standard technique)
- Alice  $\rightarrow$  Charlie: Uses random binning with artificial noise (new technique)

# Confidential MIMO BC Capacity [Liu–Liu–Poor–Shamai '10]

## Connection to [Bustin–Liu–Poor–Shamai '09]

- Generalization of the techniques of Bustin *et al.*
- Capacity region can be represented via “generalized eigenvalues”



# Confidential MIMO BC Capacity: New Derivation

- Simple derivation: Easy consequence of MIMO Wiretap treatment
- Uses standard MIMO wiretap capacity expression (no need for expression with “artificial noise”)
- Uses standard (Costa) dirty-paper coding
- Explains why dirty-paper coding is required (for at least one of the users)
- Connection to the *generalized singular values* is natural
- Allows construction of a practical scheme

## Capacity-Achieving Confidential MIMO Broadcast

$$\begin{bmatrix} \mathbf{H}_B \bar{\mathbf{K}}^{-1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix}$$

$$\begin{bmatrix} \mathbf{H}_C \bar{\mathbf{K}}^{-1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix}$$

- $\mathbf{V}_A$  of the GSVD
- Choosing directions of  $b_i > c_i$  is **optimal for Bob**
- **But...** Choosing directions of  $b_i < c_i$  is **optimal for Charlie!**



**Allocate  $b_i > c_i$  to Bob**  
**Allocate  $b_i < c_i$  to Charlie**

## Capacity-Achieving Confidential MIMO Broadcast

$$\begin{bmatrix} \mathbf{H}_B \bar{\mathbf{K}}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_B \begin{pmatrix} b_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & b_{N_A} \end{pmatrix}$$

$$\begin{bmatrix} \mathbf{H}_C \bar{\mathbf{K}}^{1/2} \mathbf{V}_A \\ \mathbf{I}_{N_A} \end{bmatrix} = \mathbf{Q}_C \begin{pmatrix} c_1 & * & * \\ 0 & \ddots & * \\ 0 & 0 & c_{N_A} \end{pmatrix}$$

## Scheme

**Charlie:** Uses SIC or DPC over elements satisfying  $c_i > b_i$

**Bob:** Uses DPC scheme over elements satisfying  $b_i > c_i$

- Alice pre-cancels interference also from Charlie's signals
- Bob can gain no info. of Charlie's message from DPC

## Capacity-Achieving Confidential MIMO Broadcast

## Decoupling the Modulation

- Alice can apply additional unitary block-matrix operations:

$$\tilde{\mathbf{T}}_B = \overbrace{\begin{pmatrix} \mathbf{Q}_{B;B}^\dagger & \mathbf{0} \\ \mathbf{0} & \mathbf{Q}_{B;C}^\dagger \end{pmatrix}}^{\text{Bob}} \begin{pmatrix} \mathbf{B}_B & * \\ \mathbf{0} & \mathbf{B}_C \end{pmatrix} \overbrace{\begin{pmatrix} \mathbf{V}_B & \mathbf{0} \\ \mathbf{0} & \mathbf{V}_C \end{pmatrix}}^{\text{Alice}}$$

$$\tilde{\mathbf{T}}_C = \overbrace{\begin{pmatrix} \mathbf{Q}_{C;B}^\dagger & \mathbf{0} \\ \mathbf{0} & \mathbf{Q}_{C;C}^\dagger \end{pmatrix}}^{\text{Charlie}} \begin{pmatrix} \mathbf{C}_B & * \\ \mathbf{0} & \mathbf{C}_C \end{pmatrix} \overbrace{\begin{pmatrix} \mathbf{V}_B & \mathbf{0} \\ \mathbf{0} & \mathbf{V}_C \end{pmatrix}}^{\text{Alice}}$$

- SVD of Charlie's "block"  $\mathbf{C}_C$  allows him to avoid SIC or DPC
- Other decompositions with desired properties can be applied (SVD of  $\mathbf{B}_C$ , 2-GMD of  $(\mathbf{B}_B, \mathbf{B}_C)$ , 2-GMD of  $(\mathbf{C}_B, \mathbf{C}_C)$ , etc.)

# Summary

- Alternative derivation of the result of Bustin *et al.* for MIMO WTC via matrix analysis tools
- Simple derivation of capacity region of confidential MIMO BC
- Connection to generalized singular values is immediate
- Explicitly shows why DPC by at least one user is required