

# Computer Arithmetic - Spring 1998

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[http://www.eng.tau.ac.il/~guy/Computer\\_Arithmetic/arith\\_home.html](http://www.eng.tau.ac.il/~guy/Computer_Arithmetic/arith_home.html)

## Assignment Number 1

**Deadline: March 19th, 16:00 (beginning of class)**

1. Prove that the sum bit  $s_n$  depends on each of the input bits.
2. Prove a lower bound on the depth of a circuit that computes a value that depends on  $n$  inputs. Assume that the fan-in of each gate is at most  $d$ . Substitute the following values in the lower bound:  $n = 64, 128$  and  $d = 2, 3, 4$ .
3. Consider the recurrence equation:

$$f(n) = \begin{cases} b & \text{if } n \leq 1 \\ a \cdot f(n/c) + bn & \text{if } n > 1 \end{cases}$$

Prove that if  $n = c^k$ , for a positive integer  $k$ , then

$$f(n) = \begin{cases} O(n) & \text{if } a < c \\ O(n \log n) & \text{if } a = c \\ O(n^{\log_c a}) & \text{if } a > c \end{cases}$$

4. Consider the recurrence equation:

$$T(n) \leq \begin{cases} 0 & \text{if } n \leq 1 \\ \min_{n_1 \leq n/2} \{T(n - n_1) + n_1\} & \text{otherwise} \end{cases}$$

Prove that  $T(n)$  is minimized if  $n_1 = n/2$  is chosen, and therefore,  $T(n) = O(n \log n)$ .

5. Prove that if  $T(MUX_n) = \log n$ , then  $T(CSA_n) = O(\log^2 n)$ .
6. Design an adder that is input two vectors  $a, b \in \{0, 1\}^n$  and computes  $s^0$  and  $s^1$  such that  $\langle s^0 \rangle = \langle a \rangle + \langle b \rangle$  and  $\langle s^1 \rangle = \langle a \rangle + \langle b \rangle + 1$ . The delay of the adder should be  $O(\log n)$  and the cost  $O(n \log n)$ .